# Towards a resource theory of contextuality



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Contextuality: a fundamental non-classical phenomenon of QM



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Contextuality as a resource for QC:

- Raussendorf (2013) MBQC
  "Contextuality in measurement-based quantum computation"
- Howard, Wallman, Veith, & Emerson (2014) MSD
  "Contextuality supplies the 'magic' for quantum computation"

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- composional aspects
- in particular, "free" operations
- A–B: qualitative hierarchy of contextuality for empirical models
- quantitative grading measure of contextuality (NB: there may be more than one useful measure)

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- Relates to quantifiable advantages in QC and QIP tasks

#### **Empirical data**



## Abramsky–Brandenburger framework

Measurement scenario  $\langle X, \mathcal{M}, O \rangle$ :

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Example: (2,2,2) Bell scenario

- The set of variables is  $X = \{a_1, a_2, b_1, b_2\}$ .
- The outcomes are  $O = \{0, 1\}$ .
- The measurement contexts are:

$$\{ \{a_1, b_1\}, \{a_1, b_2\}, \{a_2, b_1\}, \{a_2, b_2\} \}.$$

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A joint outcome or **event** in a context *C* is  $s \in O^C$ , e.g.

$$s = [a_1 \mapsto 0, b_1 \mapsto 1]$$
.

(These correspond to the cells of our probability tables.)

#### Another example: 18-vector Kochen–Specker

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- ► A set of 18 variables, X = {A,..., O}
- A set of outcomes *O* = {0, 1}
- ► A measurement cover *M* = {*C*<sub>1</sub>,..., *C*<sub>9</sub>}, whose contexts *C<sub>i</sub>* correspond to the columns in the following table:

$U_1$	$U_2$	$U_3$	$U_4$	$U_5$	$U_6$	$U_7$	$U_8$	U <sub>9</sub>
Α	A	Н	Н	В	1	Р	Р	Q
В	Ε	1	K	E	K	Q	R	R
С	F	С	G	М	Ν	D	F	М
D	G	J	L	N	0	J	L	0

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Compatibility condition: these distributions "agree on overlaps", i.e.

$$\forall_{\mathcal{C},\mathcal{C}'\in\mathcal{M}}. e_{\mathcal{C}}|_{\mathcal{C}\cap\mathcal{C}'} = e_{\mathcal{C}'}|_{\mathcal{C}\cap\mathcal{C}'}.$$

where marginalisation of distributions: if  $D \subseteq C$ ,  $d \in Prob(O^C)$ ,

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For multipartite scenarios, compatibility = the **no-signalling** principle.

A (compatible) empirical model is **non-contextual** if there exists a **global distribution**  $d \in \operatorname{Prob}(O^{\chi})$  (on the joint assignments of outcomes to all measurements) that marginalises to all the  $e_c$ :

$$\exists_{d\in \operatorname{Prob}(O^X)}, \forall_{C\in\mathcal{M}}, d|_C = e_C.$$

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#### Contextuality:

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The import of results such as Bell's and Bell–Kochen–Specker's theorems is that there are empirical models arising from quantum mechanics that are contextual.

# Strong contextuality

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E.g. K–S models, GHZ, the PR box:

А	В	(0,0)	(0,1)	(1,0)	(1,1)
$a_1$	$b_1$	$\checkmark$	×	×	$\checkmark$
$a_1$	b <sub>2</sub>	$\checkmark$	×	×	$\checkmark$
$a_2$	$b_1$	$\checkmark$	×	×	$\checkmark$
$a_2$	$b_2$	×	$\checkmark$	$\checkmark$	×



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 $\forall_{\mathcal{C}\in\mathcal{M}}. \ \mathcal{C}|_{\mathcal{C}} \leq e_{\mathcal{C}}.$ 

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Equivalently, maximum weight  $\lambda$  over all convex decompositions

$$e = \lambda e^{NC} + (1 - \lambda)e'$$

where  $e^{NC}$  is a non-contextual model.

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$$NCF(e) = \lambda$$
  $CF(e) = 1 - \lambda$ 

# (Non-)contextual fraction via linear programming

Checking contextuality of e corresponds to solving

Find 
$$\mathbf{d} \in \mathbb{R}^n$$
  
such that  $\mathbf{M}\mathbf{d} = \mathbf{v}^e$   
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Computing the non-contextual fraction corresponds to solving the following linear program:

Find	$\mathbf{c} \in \mathbb{R}^n$
maximising	1 · c
subject to	$Mc \le v^e$
and	$\mathbf{c} \geq 0$

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# E.g. Equatorial measurements on GHZ(n)



Figure: Non-contextual fraction of empirical models obtained with equatorial measurements at  $\phi_1$  and  $\phi_2$  on each qubit of  $|\psi_{\text{GHZ}(n)}\rangle$  with: (a) n = 3; (b) n = 4.

# Violations of Bell inequalities

An **inequality** for a scenario  $\langle X, \mathcal{M}, O \rangle$  is given by:

- a set of coefficients  $\alpha = \{\alpha(C, s)\}_{C \in \mathcal{M}, s \in \mathcal{E}(C)}$
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For a model *e*, the inequality reads as

$$\mathcal{B}_{lpha}(oldsymbol{e})\ \leq\ oldsymbol{R}$$
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It is called a **Bell inequality** if it is satisfied by every NC model. If it is saturated by some NC model, the Bell inequality is said to be **tight**.

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For a general (no-signalling) model e, the quantity is limited only by

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The **normalised violation** of a Bell inequality  $\langle \alpha, R \rangle$  by an empirical model *e* is the value

$$rac{\max\{0,\mathcal{B}_{lpha}(\boldsymbol{e})-\boldsymbol{R}\}}{\|lpha\|-\boldsymbol{R}}\;.$$

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- The normalised violation by e of any Bell inequality is at most CF(e).
- This is attained: there exists a Bell inequality whose normalised violation by e is exactly CF(e).
- Moreover, this Bell inequality is tight at "the" non-contextual model e<sup>NC</sup> and maximally violated by "the" strongly contextual model e<sup>SC</sup>:

$$e = \mathsf{NCF}(e)e^{\mathsf{NC}} + \mathsf{CF}(e)e^{\mathsf{SC}}$$

#### Quantifying Contextuality LP:

Find	$\mathbf{c} \in \mathbb{R}^n$
maximising	1 · c
subject to	$Mc \le v^{e}$
and	$\textbf{c} \geq \textbf{0}$

$$\boldsymbol{e} = \lambda \boldsymbol{e}^{NC} + (1 - \lambda) \boldsymbol{e}^{SC}$$
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Dual LP:

Find	$\mathbf{y} \in \mathbb{R}^m$
minimising	y ⋅ v <sup>e</sup>
subject to	$\mathbf{M}^T \mathbf{y} \ge 1$
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$$\mathbf{a} := \mathbf{1} - |\mathcal{M}|\mathbf{y}|$$



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#### Dual LP:

Find	$\mathbf{y} \in \mathbb{R}^m$
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subject to	$\mathbf{M}^T \mathbf{y} \ge 1$
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$$\bm{a}:=\bm{1}-|\mathcal{M}|\bm{y}$$

.

Find	$\mathbf{a} \in \mathbb{R}^m$
maximising	a · v <sup>e</sup>
subject to	M <sup>7</sup> a <b>≤ 0</b>
and	a ≤ 1

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subject to	${f M}{f c}\leq{f v}^{e}$
and	$\mathbf{c} \geq 0$ .

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Find $\mathbf{y} \in \mathbb{R}^m$ minimising $\mathbf{y} \cdot \mathbf{v}^e$ subject to $\mathbf{M}^T \mathbf{y} \ge \mathbf{1}$ and $\mathbf{y} \ge \mathbf{0}$ 

# $\textbf{a} := \textbf{1} - |\mathcal{M}|\textbf{y}$

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computes tight Bell inequality (separating hyperplane)

# Operations on empirical models

More than one possible measure of contextuality.

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- Monotonicity wrt operations that do not introduce contextuality
- Towards a resource theory as for entanglement (e.g. LOCC), non-locality, ...

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- These operations should not increase contextuality.
- ► Write type statements e : (X, M, O) to mean that e is a (compatible) emprical model on the scenario (X, M, O).
- The operations remind one of process algebras.


#### relabelling

 $\boldsymbol{e}: \langle \boldsymbol{X}, \mathcal{M}, \boldsymbol{\mathcal{O}} \rangle, \ \alpha: (\boldsymbol{X}, \mathcal{M}) \cong (\boldsymbol{X}', \boldsymbol{M}') \ \rightsquigarrow \ \boldsymbol{e}[\alpha]: \langle \boldsymbol{X}', \mathcal{M}', \boldsymbol{\mathcal{O}} \rangle$ 

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For  $C \in \mathcal{M}, s : \alpha(C) \longrightarrow O, e[\alpha]_{\alpha(C)}(s) := e_{C}(s \circ \alpha^{-1})$ 

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$$\begin{array}{l} \text{For } C' \in M', s : C' \longrightarrow O, \, (e \upharpoonright \mathcal{M}')_{C'}(s) := e_C|_{C'}(s) \\ \quad \text{with any } C \in \mathcal{M} \text{ s.t. } C' \subseteq C \end{array}$$

#### relabelling

 $\boldsymbol{e}:\langle \boldsymbol{X}, \mathcal{M}, \boldsymbol{O} \rangle, \; \alpha: (\boldsymbol{X}, \mathcal{M}) \cong (\boldsymbol{X}', \boldsymbol{M}') \; \rightsquigarrow \; \boldsymbol{e}[\alpha]: \langle \boldsymbol{X}', \mathcal{M}', \boldsymbol{O} \rangle$ 

For  $C \in \mathcal{M}, s : \alpha(C) \longrightarrow O, e[\alpha]_{\alpha(C)}(s) := e_C(s \circ \alpha^{-1})$ 

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#### coarse-graining

$$e: \langle X, \mathcal{M}, O \rangle, \ f: O \longrightarrow O' \ \rightsquigarrow \ e/f: \langle X, \mathcal{M}, O' \rangle$$

For 
$$C \in M, s : C \longrightarrow O', (e/f)_C(s) := \sum_{t:C \longrightarrow O, f \circ t = s} e_C(t)$$



mixing

 $\boldsymbol{e}:\langle \boldsymbol{X}, \mathcal{M}, \boldsymbol{O} \rangle, \ \boldsymbol{e}':\langle \boldsymbol{X}, \mathcal{M}, \boldsymbol{O} \rangle, \lambda \in [0, 1] \ \rightsquigarrow \ \boldsymbol{e} +_{\lambda} \ \boldsymbol{e}':\langle \boldsymbol{X}, \mathcal{M}, \boldsymbol{O} \rangle$ 

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For  $C \in M$ ,  $(e\&e')_C := e_C$ For  $D \in M'$ ,  $(e\&e')_D := e'_D$ 

#### tensor

 $\textit{e}: \langle \textit{X}, \mathcal{M}, \textit{O} \rangle, \textit{ e'}: \langle \textit{X'}, \mathcal{M'}, \textit{O} \rangle \iff \textit{e} \otimes \textit{e'}: \langle \textit{X} \sqcup \textit{X'}, \mathcal{M} \star \mathcal{M'}, \textit{O} \rangle$ 

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$$\mathcal{M} \star \mathcal{M}' := \{ \mathcal{C} \sqcup \mathcal{D} \mid \mathcal{C} \in \mathcal{M}, \mathcal{D} \in \mathcal{M}' \}$$

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$$\begin{split} \mathcal{M} \star \mathcal{M}' &:= \{ C \sqcup D \mid C \in \mathcal{M}, D \in \mathcal{M}' \} \\ \mathsf{For} \; C \in \mathcal{M}, D \in \mathcal{M}', s = \langle s_1, s_2 \rangle : C \sqcup D \longrightarrow O, \\ & (e \otimes e')_{C \sqcup D} \langle s_1, s_2 \rangle := e_C(s_1) \, e'_D(s_2) \end{split}$$

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tensor

 $\begin{aligned} \mathsf{CF}(e_1 \otimes e_2) &= \mathsf{CF}(e_1) + \mathsf{CF}(e_2) - \mathsf{CF}(e_1)\mathsf{CF}(e_2) \\ \mathsf{NCF}(e_1 \otimes e_2) &= \mathsf{NCF}(e_1)\mathsf{NCF}(e_2) \end{aligned}$ 

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- ► Measure of contextuality ~→ to quantify such advantages.

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- ► Raussendorf (2013): If an ℓ2-MBQC deterministically computes a non-linear Boolean function f : 2<sup>m</sup> → 2<sup>l</sup> then the resource must be strongly contextual.
- Probabilistic version: non-linear function computed with sufficently large probability of success implies contextuality.

### average distance between two Boolean functions

$$f,g: 2^m \longrightarrow 2^r:$$
  
 $\widetilde{d}(f,g):=2^{-m}|\{\mathbf{i}\in 2^m \mid f(\mathbf{i})\neq g(\mathbf{i})\}$ 

- ► average distance between two Boolean functions  $f, g : 2^m \longrightarrow 2^l$ :  $\tilde{d}(f,g) := 2^{-m} | \{ i \in 2^m | f(i) \neq g(i) \}$
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- Then,  $1 \bar{p}_S \ge \text{NCF}(e)\tilde{\nu}(f)$ .

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▶ 
$$1 - \bar{p}_S \leq \operatorname{NCF}(e) \frac{(n-k)}{n}$$
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- Negative Probabilities Measure
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  - What (else) is this resource useful for?



# ?