#### Random Optimisation Problems

Amin Coja-Oghlan

partly based on joint work with Philipp Loick, Balazs Mezei, Greg Sorkin *thanks to Philipp for helping with the pictures!*

Goethe Universität Frankfurt

## Part 1

#### **Motivation**

- $\triangleright$  what are random optimisation problems?
- $\blacktriangleright$  and what are they good for?

## What are random optimisation problems?



#### Examples

- $\triangleright$  minimum spanning trees with random weights [F85,...]
- planted clique [AKS98,...]
- $\triangleright$  MAX CUT on random graphs [CGHS04,...]
- 



Hypothesis: hard instances are rare. . .

- $\triangleright$  ... but real-world instances are easy!
- random models of real-world problems?
- ▶ semi-random, smoothed, power law... [FK01,ST01,FR19]
- $\triangleright$  where are the really hard problems?



#### Number of solved instances CPU time 4,000s

#### Example: SAT solving

- $\blacktriangleright$  10m variables, 32m clauses
- ▶ 2020 winner KISSAT based on clause learning [BFFH2020]
- $\triangleright$  real-world SAT instances solved on an industrial scale



#### Number of solved instances CPU time 4,000s

#### Example: random *k*-SAT

- <sup>Ï</sup> satisfiability threshold *m*/*n* ∼ 2 *k*
- ► algorithmic barrier  $m/n \sim 2^k$

 $[DSS15]$ log(*k*)/*k* [ACO08,CO09]





Hypothesis: hard instances are rare. . .

- ► ... but real-world instances are easy! *[they are structured]*
- **Example 1** can random instances model real-world problems? *[not really]*
- ▶ semi-random, smoothed, power law. . *[excess entropy]*
- where are the really hard problems? *<i>[try random ones!]*

## The probabilistic method



Success stories [personal selection]

- low-density parity check codes [G63,KRU11]
- compressed sensing [KMSSZ12,DJM13]
- 

► group testing [COGHKL20]

# A glimpse of complexity?

- ▶ proof complexity [BSW00,...]
- 
- 
- ▶ overlap gap property [GS14,GZ19,BAWZ20,...]
- statistical algorithms [FPV13,SW20]

• planted clique vs SOS [DM15,BHKKMP19,...] ► *k*-SAT refutation [F02,FO07,COGL04,KMODW17,...]

#### **Techniques**

- $\triangleright$  what can we prove about random optimisation problems?
- $\triangleright$  what techniques do we have at our disposal?

## Running example



.

#### MAXCUT on random regular graphs

 $G = G(n, d)$  random *d*-regular graph of order *n* 

$$
\text{MAXCUT}(\mathbb{G}) = \max_{\sigma \in \{\pm 1\}^n} \sum_{v w \in E(\mathbb{G})} \frac{1 - \sigma_v \sigma_w}{2}
$$

## Combinatorial bounds



Greedy algorithms [DDSW03]

- $\triangleright$  assign each vertex the minority spin amongst its neighbours
- $\blacktriangleright$  method of differential equations

## Combinatorial bounds



#### The first moment bound

The expected number of cuts of size *αdn*/2 equals

$$
\exp\left(n\left[\left(1-\frac{d}{2}\right)\log 2+\frac{d}{2}H(\alpha)+o(1)\right]\right)
$$

## The physics perspective



#### The Ising antiferromagnet

For an inverse temperature  $\beta \geq 0$  introduce

$$
\mathcal{H}_{\mathbb{G}}(\sigma) = \sum_{vw \in E(\mathbb{G})} \frac{1 + \sigma_v \sigma_w}{2}
$$

$$
Z_{\mathbb{G},\beta} = \sum_{\sigma \in \{\pm 1\}^n} \exp(-\beta \mathcal{H}_{\mathbb{G}}(\sigma))
$$

$$
\mu_{\mathbb{G},\beta}(\sigma) = \exp(-\beta \mathcal{H}_{\mathbb{G}}(\sigma)) / Z_{\mathbb{G},\beta}
$$

# The physics perspective



#### The Ising antiferromagnet

 $\blacktriangleright$  The average number of edges cut equals

$$
\frac{dn}{2}-\langle \mathcal{H}_{\mathbb{G}}(\boldsymbol{\sigma}),\mu_{\mathbb{G},\beta}\rangle=\frac{dn}{2}+\frac{\partial}{\partial\beta}\log Z_{\mathbb{G},\beta}
$$

 $\triangleright$  Taking *β* → ∞ yields

$$
\text{MAXCUT}(\mathbb{G}) = \lim_{\beta \to \infty} \frac{dn}{2} + \frac{\partial}{\partial \beta} \log Z_{\mathbb{G}, \beta}
$$



#### Absence of long-range correlations [KMRTSZ07]

Gibbs uniqueness: the root decouples from the boundary Non-reconstruction: decoupling from a typical boundary Static replica symmetry:  $\mathbb{E} \left| \mu_{\mathbb{G},\beta} \left( \{ \boldsymbol{\sigma}_{v_1} = \boldsymbol{\sigma}_{v_2} \} \right) - \frac{1}{2} \right|$  $\frac{1}{2}$  =  $o(1)$ 

# Replica symmetry breaking



#### First moment redux

- $\triangleright$  assume static replica symmetry, assemble G one edge at a time
- $\blacktriangleright$  every time we insert a new edge  $e = vw$ ,

$$
\frac{Z_{\mathbb{G},\beta}(\mathbb{G} + e)}{Z_{\mathbb{G},\beta}} \sim 1 - (1 - e^{-\beta})\mu_{\mathbb{G},\beta}(\{\sigma_v = \sigma_w\}) \sim \frac{1 + e^{-\beta}}{2}
$$

$$
\frac{1}{n}\log Z_{\mathbb{G},\beta} \sim \log 2 + \frac{d}{2}\log\frac{1 + e^{-\beta}}{2} = \frac{1}{n}\log \mathbb{E}[Z_{\mathbb{G},\beta}]
$$

## Replica symmetry breaking



#### Pure states [KMRTSZ07]

low  $β$  the cross-section is contiguous high  $\beta$  decomposition into separate pure states



#### The overlap [KMRTSZ07]

replica symmetry the overlap  $\sigma \cdot \sigma'/n$  concentrates about 0 1-step replica symmetry breaking concentration on two points full replica symmetry breaking no overlap concentration

The replica symmetry breaking transition

The free energy The singularities of

$$
\phi_d(\beta) = \lim_{n \to \infty} \frac{1}{n} \mathbb{E}[\log Z_{\mathbb{G},\beta}]
$$

are called *phase transitions*.

## The replica symmetry breaking transition



Theorem [COLMS20]

For any  $d \geq 3$  we have

$$
\phi_d(\beta) = \log 2 + \frac{d}{2} \log \frac{1 + e^{-\beta}}{2} \quad \text{if } \beta \le \beta^*(d) = \log \frac{\sqrt{d-1} + 1}{\sqrt{d-1} - 1}
$$

$$
\phi_d(\beta) < \log 2 + \frac{d}{2} \log \frac{1 + e^{-\beta}}{2} \quad \text{if } \beta > \beta^*(d)
$$

## The replica symmetry breaking transition



Non-reconstruction/Kesten-Stigum

- ► broadcasting on the  $(d-1)$ -ary tree
- ► reconstruction threshold at  $\beta^*(d) = \log \frac{\sqrt{d-1}+1}{\sqrt{d-1}-1}$
- $\blacktriangleright$  Gibbs uniqueness threshold at  $β<sup>†</sup>(d) = log \frac{d}{d-2}$

## First attempt

#### The second moment method

$$
\mathbb{E}[Z_{\mathbb{G},\beta}^{2}] = \sum_{\sigma,\sigma'\in\{\pm1\}^{n}} \mathbb{E}\left[\exp(-\beta\mathcal{H}_{\mathbb{G}}(\sigma) - \beta\mathcal{H}_{\mathbb{G}}(\sigma'))\right]
$$

$$
= \sum_{a=-n}^{n} \sum_{\sigma,\sigma'=a} \mathbb{E}\left[\exp(-\beta\mathcal{H}_{\mathbb{G}}(\sigma) - \beta\mathcal{H}_{\mathbb{G}}(\sigma'))\right]
$$

$$
= \exp\left(n \max_{\alpha\in[-1,1]} f_{d,\beta}(\alpha) + o(n)\right)
$$

$$
f_{d,\beta}(\alpha) = (1-d)\log 2 + H\left(\frac{1+\alpha}{2}\right) + \frac{d}{2}\log\left((1+e^{-\beta})^2 + \alpha^2(1-e^{-\beta})^2\right).
$$

### First attempt





#### The second moment method

- $▶ E[Z_{\text{G},\beta}^2] = O(\mathbb{E}[Z_{\text{G},\beta}]^2)$  iff  $f_{d,\beta}(\alpha)$  attains its max at  $\alpha = 0$
- $\blacktriangleright$  this is the case iff *β* < *β*<sup>\*</sup>(*d* + 1) < *β*<sup>\*</sup>(*d*)

### First attempt



#### The Erdős-Rényi model

- $\rightarrow$  identical first/second moments
- $\blacktriangleright$  Po(*d*) offspring in the Erdős-Rényi model
- lacktriangleright phase transition at  $\beta^*$

[MNS15]

#### . . . aka stochastic block model

- ► draw a configuration  $\boldsymbol{\sigma}^* \in \{\pm 1\}^n$
- ► draw a *d*-regular graph G<sup>\*</sup> from

$$
\mathbb{P}\left[\mathbb{G}^* = G \,|\, \boldsymbol{\sigma}^*\right] \propto \exp(-\beta \mathcal{H}_G(\boldsymbol{\sigma}^*))
$$

$$
\blacktriangleright \mathbb{P}[\mathbb{G}^* \in \mathscr{E}] = \Theta(\mathbb{E}[Z_{\mathbb{G},\beta} \cdot \mathbb{I}\{\mathbb{G} \in \mathscr{E}\}])
$$



#### . . . aka stochastic block model

- ► draw a configuration  $\boldsymbol{\sigma}^* \in \{\pm 1\}^n$
- ► draw a *d*-regular graph G<sup>\*</sup> from

$$
\mathbb{P}\left[\mathbb{G}^* = G \,|\, \boldsymbol{\sigma}^*\right] \propto \exp(-\beta \mathcal{H}_G(\boldsymbol{\sigma}^*))
$$

$$
\blacktriangleright \mathbb{P}[\mathbb{G}^* \in \mathscr{E}] = \Theta(\mathbb{E}[Z_{\mathbb{G},\beta} \cdot \mathbb{I}\{\mathbb{G} \in \mathscr{E}\}])
$$



#### The broadcasting process

- ► locally ( $\mathbb{G}^*$ , *σ*<sup>\*</sup>) resembles the broadcasing process
- $\blacktriangleright \implies \mu_{\mathbb{G}^*,\beta}$  has static replica symmetry for  $\beta < \beta^*(d)$
- $\blacktriangleright$   $\Rightarrow$  overlap concentrates about zero for  $\beta < \beta^*$ (*d*) [cf. MNS15]

 $\mu_{\mathbb{G}^*,\beta}(\{\boldsymbol{\sigma}\cdot\boldsymbol{\sigma}'/n<\varepsilon\})\sim 1$ 

#### The truncated second moment

- $\mathcal{E} = {\mu_{\mathbb{G},\beta}(\{\boldsymbol{\sigma}\cdot\boldsymbol{\sigma}'/n<\varepsilon\}) \sim 1\}}$
- $\blacktriangleright$  E[ $Z_{\mathbb{G},\beta}$  · 1{ $\mathbb{G}\in\mathscr{E}$ }] ~ E[ $Z_{\mathbb{G},\beta}$ ]
- $\blacktriangleright$   $\mathbb{E}[Z_{\mathbb{G},\beta}^2 \cdot \mathbb{I}\{\mathbb{G} \in \mathscr{E}\}] = \mathbb{E}[Z_{\mathbb{G},\beta}]^2 \cdot \exp(\varepsilon^2 n)$
- $\blacktriangleright$  Paley-Zygmund inequality

⇒  $Z_{\mathbb{G},\beta}$  ≥  $Z_{\mathbb{G},\beta}$  · 1{ $\mathbb{G} \in \mathscr{E}$ } ≥  $\mathbb{E}[Z_{\mathbb{G},\beta}] \cdot \exp(o(n))$  w.h.p.

*The first part of the theorem follows.*

## The planted model diverges

#### Quiet vs noisy planting

- $\triangleright$  log *Z*<sub>G,*β*</sub> is tightly concetrated about  $\mathbb{E}[\log Z_{\text{G},\beta}]$
- $\blacktriangleright$   $\Rightarrow$  if  $\mathbb{E}[\log Z_{\mathbb{G},\beta}]$  ~  $\log \mathbb{E}[Z_{\mathbb{G},\beta}]$ , then  $\mathbb{G} \approx \mathbb{G}^*$

$$
\lim_{n \to \infty} \frac{1}{n} \mathbb{E}[\log Z_{\mathbb{G}^*,\beta}] > \log 2 + \frac{d}{2} \log \frac{1 + e^{-\beta}}{2} \ge \phi_d(\beta)
$$
  
\n
$$
\Rightarrow \phi_d(\beta) < \log 2 + \frac{d}{2} \log \frac{1 + e^{-\beta}}{2}
$$

## The planted model diverges

The planted free energy [COHKLMPP20]

$$
\lim_{n \to \infty} \frac{1}{n} \mathbb{E}[\log Z_{\mathbb{G}^*,\beta}] = \sup_{\pi} \mathcal{B}_{d,\beta}^*(\pi)
$$

$$
\mathcal{B}_{d,\beta}^*(\pi) = \mathbb{E} \Big[ \frac{\Lambda \left( \sum_{\sigma = \pm 1} \prod_{i=1}^d 1 - (1 - e^{-\beta})(1 + \sigma \rho_i)/2 \right)}{2^{1 - d}(1 + e^{-\beta})^d}
$$

$$
- \frac{d\Lambda \left( 1 - (1 - e^{-\beta})(1 + \rho_1 \rho_2)/2 \right)}{1 + e^{-\beta}} \Big]
$$

 $Λ(x) = x \log x$ 

## The planted model diverges

#### Lemma

For  $\beta > \beta^*(d)$  there is  $\varepsilon > 0$  such that  $\pi_{\varepsilon} = \frac{1}{2}$  $\frac{1}{2}$  ( $\delta$ <sub>−*ε*</sub> +  $\delta$ *<sub>ε</sub>*) satisfies

$$
\mathcal{B}_{d,\beta}^*(\pi_\varepsilon) > \log 2 + \frac{d}{2} \log \frac{1 + e^{-\beta}}{2}
$$

#### Proof

- ► Compute the first four derivatives of  $\varepsilon \mapsto \mathcal{B}^*_{d,\beta}(\pi_{\varepsilon})$
- $\blacktriangleright$  (seriously...)

*We thus obtain the second part of the theorem.*

#### Theorem [SSZ16]



For any  $d \ge 3$ ,  $\beta > 0$ , any probability measure *r* on (−1, 1) and any  $0 < y < 1$  we have

$$
\phi_d(\beta) \le \frac{1}{y} \log \mathbb{E}[X^y] - \frac{d}{2y} \log \mathbb{E}[Y^y]
$$
  

$$
X = \sum_{\tau = \pm 1} \prod_{i=1}^d 1 - (1 - e^{-\beta}) \frac{1 + \tau \rho_i}{2}
$$
  

$$
Y = 1 - (1 - e^{-\beta}) \sum_{\tau = \pm 1} \frac{(1 + \tau \rho_1)(1 + \tau \rho_2)}{4}
$$

*Regular variant of [G03, FL03, PT04].*





Pure states and Poisson–Dirichlet weights

- $\blacktriangleright$  auxiliary model that represents the 1-step rsb scenario
- $\rightarrow$  interpolation along time *t* ∈ [0, 1]
- ▶ *Hypothesis:* the weights of the pure states are Poisson–Dirichlet
- $\triangleright$  to bound MAX CUT we use MAXCUT(<del>G</del>) ≤  $\frac{dn}{2}$  $\frac{\ln}{2} + \frac{1}{\beta}$ *β* log*Z*G,*β*

Theorem

For any  $d \geq 3$  we have

$$
\text{MAXCUT}(\mathbb{G}) \le \frac{dn}{2} \left( 1 + \inf_{\alpha, z} -\frac{\log \zeta \mathscr{A}^d \zeta}{\log z} + \frac{d \log(1 - 2\alpha^2 + 2\alpha^2 z)}{2 \log z} \right)
$$
  

$$
\zeta = (1, 0, 0, \dots) \qquad \zeta = (1, z^{-1/2}, z^{-1}, \dots)^T
$$
  

$$
\mathscr{A} = (1 - 2\alpha) \text{id} + 2\alpha \sqrt{z} \begin{bmatrix} 0 & 1 & 0 & \cdots & \cdots & \cdots \\ \frac{1}{2} & 0 & \frac{1}{2} & \cdots & \cdots & \cdots \\ 0 & \frac{1}{2} & 0 & \frac{1}{2} & \cdots & \cdots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots \\ \vdots & \ddots & \ddots & \ddots & \ddots & \ddots \\ \end{bmatrix}
$$



#### Some numbers



*Match the numbers of [ZB10].*

# **Outlook**



## Summary

#### Random optimisaton problems

- $\blacktriangleright$  hard but rewarding
- $\blacktriangleright$  algorithmic challenges and new algorithms
- $\blacktriangleright$  probabilistic constructions
- $\blacktriangleright$  techniques:
	- **•** physics-enhanced moment methods
	- $\rightarrow$  the planted model
	- $\rightarrow$  spatial mixing and asymptotic Gibbs measures
	- $\rightarrow$  the interpolation method
	- $\rightarrow$  coupling arguments