

A Theory of Trotter Error

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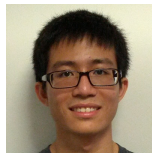
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[arXiv:1901.00564/PRL](https://arxiv.org/abs/1901.00564)

[arXiv:1912.08854](https://arxiv.org/abs/1912.08854)

Quantum simulation

- Dynamics of a quantum system are given by its Hamiltonian $\mathcal{H}(t)$ according to the Schrödinger equation

$$\frac{d}{dt}\mathcal{U}(t) = -i\mathcal{H}(t)\mathcal{U}(t), \quad \mathcal{U}(0) = I.$$

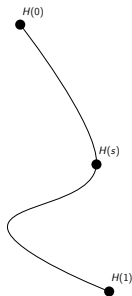
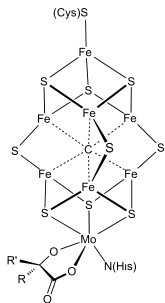
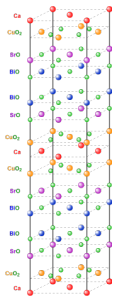
- We formally write the solution $\mathcal{U}(t) = \exp_{\mathcal{T}}\left(-i\int_0^t d\tau\mathcal{H}(\tau)\right)$. When $\mathcal{H}(t) \equiv H$ is time-independent, we have closed-form solution $\mathcal{U}(t) = e^{-itH}$.

Quantum simulation problem

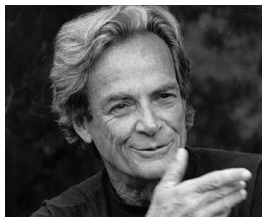
Given a description of the Hamiltonian H and evolution time t , perform e^{-itH} up to some error ϵ (in spectral norm):

$$\|U - e^{-itH}\| \leq \epsilon.$$

Reasons to study quantum simulation



$$\boxed{A} \boxed{x} = \boxed{b}$$



“...nature isn't classical, dammit, and if you want to make a simulation of nature, you'd better make it quantum mechanical, and by golly it's a wonderful problem, because it doesn't look so easy.”

— Richard Feynman

Product formulas

- Also known as Trotterization or the splitting method.
- Target system: $H = \sum_{\gamma=1}^r H_{\gamma}$, where each H_{γ} is Hermitian and can be exponentiated with cost $\mathcal{O}(1)$.
- Can use the first-order Lie-Trotter formula¹

$$\mathcal{S}_1(t) := e^{-itH_r} \dots e^{-itH_1} = e^{-itH} + \mathcal{O}(t^2)$$

with **Trotter error** $\mathcal{O}(t^2)$.

- To simulate for a large t , divide the evolution into r **Trotter steps** and simulate each step with error at most ϵ/r .
- Choose the **Trotter number** r to be sufficiently large so that the entire simulation has error at most ϵ .

¹[Lloyd 96]

Higher-order product formulas

- A general p th-order product formula takes the form

$$\mathcal{S}_p(t) := \prod_{v=1}^{\Upsilon} \prod_{\gamma=1}^{\Gamma} e^{-ita_{(v,\gamma)} H_{\pi_v(\gamma)}} = e^{-itH} + O(t^{p+1}).$$

Higher-order Suzuki formulas²

The $(2k)$ th-order Suzuki formula $\mathcal{S}_{2k}(t) = e^{-itH} + O(t^{2k+1})$ is defined recursively by

$$\begin{aligned}\mathcal{S}_2(t) &:= e^{-i\frac{t}{2}H_1} \dots e^{-i\frac{t}{2}H_r} e^{-i\frac{t}{2}H_r} \dots e^{-i\frac{t}{2}H_1}, \\ \mathcal{S}_{2k}(t) &:= \mathcal{S}_{2k-2}(u_k t)^2 \mathcal{S}_{2k-2}((1 - 4u_k)t) \mathcal{S}_{2k-2}(u_k t)^2,\end{aligned}$$

where $u_k := 1/(4 - 4^{1/(2k-1)})$.

²[Suzuki 92]

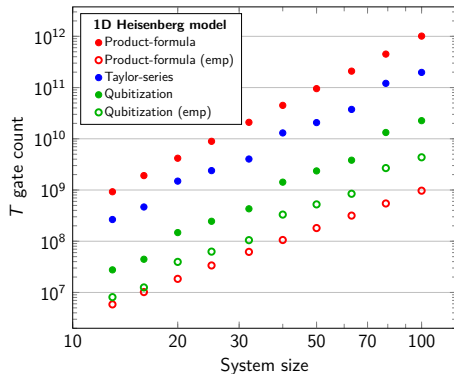
Other simulation algorithms

- Recent algorithms have improved asymptotic performance as a function of t and ϵ over the product-formula approach...



Other simulation algorithms

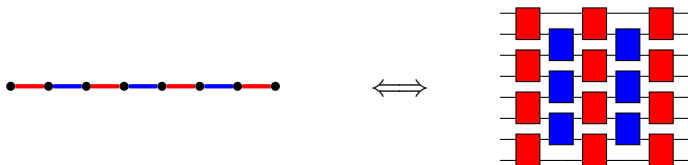
- Recent algorithms have improved asymptotic performance as a function of t and ϵ over the product-formula approach...
- ... but the empirical performance of product formulas can be significantly better.³



³[Childs, Maslov, Nam, Ross, Su 18]

Reasons to study product formulas

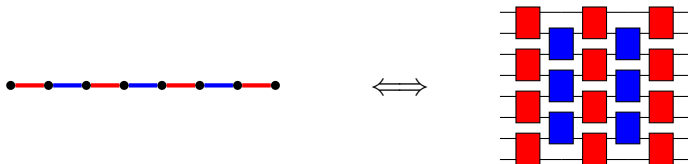
- The product-formula algorithm is ancilla-free and is arguably the simplest approach to quantum simulation.
- Product formulas can use Hamiltonian commutativity to give surprisingly efficient simulation in practice.
- Product formulas can preserve locality of the simulated system, which can be used to reduce the simulation cost.



- Other applications: classical simulation of quantum systems, numerical analysis...

Reasons to study product formulas

- The product-formula algorithm is ancilla-free and is arguably the simplest approach to quantum simulation.
- **Product formulas can use Hamiltonian commutativity to give surprisingly efficient simulation in practice.**
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Previous analyses of Trotter error

- For sufficiently small t , Trotter error can be represented exactly by the Baker–Campbell–Hausdorff formula:

$$e^{-itB}e^{-itA} = e^{-it(A+B) - \frac{t^2}{2}[B,A] + i\frac{t^3}{12}[B,[B,A]] - i\frac{t^3}{12}[A,[B,A]] + \dots}$$

- Truncating the BCH expansion ignores significant, potentially dominant contributions of Trotter error.⁴
- Using tail bounds does not exploit the commutativity of Hamiltonian summands.⁵
- Infinite-series expansion is only advantageous for systems with Lie-algebraic structure.⁶

⁴[Wecker, Bauer, Clark, Hastings, Troyer 14]

⁵[Berry, Ahokas, Cleve, Sanders 07], [Bravyi, Gosset 17]

⁶[Somma 16]

Analysis of the first-order formula

- For Hamiltonian $H = A + B$ and $t \geq 0$, the first-order formula $\mathcal{S}_1(t) = e^{-itB}e^{-itA}$ satisfies the differential equation

$$\frac{d}{dt}\mathcal{S}_1(t) = -iH\mathcal{S}_1(t) + e^{-itB}\left(e^{itB}iAe^{-itB} - iA\right)e^{-itA}$$

with initial condition $\mathcal{S}_1(0) = I$.

- Using the variation-of-parameters formula,

$$\mathcal{S}_1(t) = e^{-itH} + \int_0^t d\tau_1 e^{-i(t-\tau_1)H} e^{-i\tau_1 B} \left(e^{i\tau_1 B} iAe^{-i\tau_1 B} - iA \right) e^{-i\tau_1 A}.$$

- We further have

$$e^{i\tau_1 B} iAe^{-i\tau_1 B} - iA = \int_0^{\tau_1} d\tau_2 e^{i\tau_2 B} [iB, iA] e^{-i\tau_2 B}.$$

Analysis of the first-order formula

- Altogether, we have the integral representation

$$\mathcal{S}_1(t) - e^{-itH} = \int_0^t d\tau_1 \int_0^{\tau_1} d\tau_2 e^{-i(t-\tau_1)H} e^{-i\tau_1 B} e^{i\tau_2 B} [iB, iA] e^{-i\tau_2 B} e^{-i\tau_1 A}$$

and the error bound

$$\|\mathcal{S}_1(t) - e^{-itH}\| \leq \frac{t^2}{2} \|[B, A]\|.$$

A multi-term Hamiltonian $H = \sum_{\gamma=1}^{\Gamma} H_{\gamma}$ can be handled by bootstrapping the above bound.

- A similar bound holds for the second-order formula.
- Generalization to arbitrary higher-order formulas was previously unknown.

Trotter error with commutator scaling

Trotter error with commutator scaling

A p th-order product formula $\mathcal{S}_p(t)$ can approximate the evolution of Hermitian $H = \sum_{\gamma=1}^{\Gamma} H_{\gamma}$ for $t \geq 0$ with Trotter error

$$\left\| \mathcal{S}_p(t) - e^{-itH} \right\| = \mathcal{O}\left(\tilde{\alpha}_{\text{comm}} t^{p+1}\right),$$

where $\tilde{\alpha}_{\text{comm}} := \sum_{\gamma_1, \gamma_2, \dots, \gamma_{p+1}} \left\| \left[H_{\gamma_{p+1}}, \dots \left[H_{\gamma_2}, H_{\gamma_1} \right] \right] \right\|$.

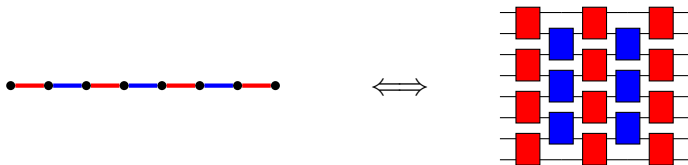
- Need $\mathcal{O}\left(\Gamma \tilde{\alpha}_{\text{comm}}^{1/p} t^{1+1/p}\right)$ gates to achieve constant accuracy.
- Asymptotic complexity is independent of how the Hamiltonian summands are ordered.
- Related bounds exist for multiplicative error and imaginary-time evolution.

Nearest-neighbor lattice Hamiltonian

1D Nearest-neighbor lattice Hamiltonian

$$H = \sum_{j=1}^{n-1} H_{j,j+1}, \text{ where } H_{j,j+1} \text{ acts only on qubits } j \text{ and } j+1.$$

- Models many important physical systems in condensed matter physics, nuclear physics, and quantum field theory.
- Previously claimed without rigorous justification that product formulas have gate complexity $(nt)^{1+o(1)}$.⁷



⁷[Jordan, Lee, Preskill 12]

Nearest-neighbor lattice Hamiltonian

- For nearest-neighbor interactions, we can use locality to simplify $\tilde{\alpha}_{\text{comm}}$, giving

$$\tilde{\alpha}_{\text{comm}} = \sum_{\gamma_1, \gamma_2, \dots, \gamma_{p+1}} \left\| \left[H_{\gamma_{p+1}}, \dots, \left[H_{\gamma_2}, H_{\gamma_1} \right] \right] \right\| = \mathcal{O}(n).$$

- **We thus proved the Jordan-Lee-Preskill claim**, giving a simple lattice simulation with nearly optimal gate complexity.⁸

⁸[Haah, Hastings, Kothari, Low 18]

Electronic structure Hamiltonian

Second-quantized plane-wave electronic structure

$$H = \underbrace{\frac{1}{2n} \sum_{j,k,\nu} \kappa_\nu^2 \cos[\kappa_\nu \cdot r_{k-j}] A_j^\dagger A_k}_{T} - \underbrace{\frac{4\pi}{\omega} \sum_{j,l,\nu \neq 0} \frac{\zeta_l \cos[\kappa_\nu \cdot (\tilde{r}_l - r_j)]}{\kappa_\nu^2} N_j}_{U} + \underbrace{\frac{2\pi}{\omega} \sum_{\substack{j \neq k \\ \nu \neq 0}} \frac{\cos[\kappa_\nu \cdot r_{j-k}]}{\kappa_\nu^2} N_j N_k}_{V}.$$

- An efficient simulation can help design and engineer new pharmaceuticals, catalysts and materials.
- To simulate n spin orbitals for time t , the best previous approach is the interaction-picture method⁹ with complexity $\tilde{O}(n^2 t)$ (and a likely large prefactor).

⁹[Low, Wiebe 18]

Electronic structure Hamiltonian

- Using fermionic Fourier transform, we can diagonalize the kinetic term in the plane-wave basis

$$\frac{1}{2n} \sum_{j,k,\nu} \kappa_\nu^2 \cos[\kappa_\nu \cdot r_{k-j}] A_j^\dagger A_k = \text{FFFT}^\dagger \left(\frac{1}{2} \sum_\nu \kappa_\nu^2 N_\nu \right) \text{FFFT}.$$

- Using the commutation rules of second-quantized operators, we estimate

$$\tilde{\alpha}_{\text{comm}} = \sum_{\gamma_1, \gamma_2, \dots, \gamma_{p+1}} \left\| \left[H_{\gamma_{p+1}}, \dots, [H_{\gamma_2}, H_{\gamma_1}] \right] \right\| = \mathcal{O}(n^{p+1}).$$

- We thus showed that product formulas have gate complexity $n^{2+o(1)} t^{1+o(1)}$, confirming a recent numerical study.¹⁰**

¹⁰[Kivlichan, Gidney, Berry et al. 19]

k -local Hamiltonian

k -local Hamiltonian

$H = \sum_{j_1, \dots, j_k} H_{j_1, \dots, j_k}$, where each H_{j_1, \dots, j_k} acts only on $k = \mathcal{O}(1)$ qubits j_1, \dots, j_k .

- Ubiquitous in physics.
- The best previous algorithm is the qubitization approach¹¹ with complexity $\tilde{\mathcal{O}}(n^k \|H\|_1 t)$, scaling with the 1-norm

$$\|H\|_1 := \sum_{j_1, \dots, j_k} \|H_{j_1, \dots, j_k}\|.$$

¹¹[Low, Chuang 19]

k -local Hamiltonian

- Using locality to simplify $\tilde{\alpha}_{\text{comm}}$, we obtain

$$\tilde{\alpha}_{\text{comm}} = \sum_{\gamma_1, \gamma_2, \dots, \gamma_{p+1}} \left\| \left[H_{\gamma_{p+1}}, \dots, [H_{\gamma_2}, H_{\gamma_1}] \right] \right\| = \mathcal{O}(\|H\|_1^p \|H\|_1),$$

where

$$\|H\|_1 := \max_l \max_{j_l} \sum_{j_1, \dots, j_{l-1}, j_{l+1}, \dots, j_k} \|H_{j_1, \dots, j_k}\|.$$

- We thus gave a simulation algorithm with complexity $n^k \|H\|_1 \|H\|_1^{o(1)} t^{1+o(1)}$ scaling with the induced 1-norm.
- We have the norm inequality

$$\|H\|_1 := \max_l \max_{j_l} \sum_{j_1, \dots, j_{l-1}, j_{l+1}, \dots, j_k} \|H_{j_1, \dots, j_k}\| \leq \|H\|_1 := \sum_{j_1, \dots, j_k} \|H_{j_1, \dots, j_k}\|$$

and the gap can be significant for many k -local Hamiltonians.

Rapidly decaying power-law Hamiltonian

Rapidly decaying power-law Hamiltonian

$H = \sum_{\vec{i}, \vec{j}} H_{\vec{i}, \vec{j}}$, where $H_{\vec{i}, \vec{j}}$ acts only on qubits $\vec{i}, \vec{j} \in \Lambda \subset \mathbb{R}^d$ and

$$\|H_{\vec{i}, \vec{j}}\| \leq \begin{cases} 1, & \text{if } \vec{i} = \vec{j}, \\ \frac{1}{\|\vec{i} - \vec{j}\|_2^\alpha}, & \text{if } \vec{i} \neq \vec{j}, \end{cases}$$

with dimension d and exponent $\alpha > 2d$.

- Examples include the dipole-dipole interactions ($\alpha = 3$) and the Van der Waals interactions ($\alpha = 6$).
- The best previous approach is an algorithm based on Lieb-Robinson bounds¹² using $\tilde{O}((nt)^{1+2d/(\alpha-d)})$ gates.

¹²[Tran, Guo et al. 19]

Rapidly decaying power-law Hamiltonian

- We truncate terms with distance larger than a cut-off $\ell = \Theta\left((nt/\epsilon)^{1/(\alpha-d)}\right)$.
- For the truncated Hamiltonian, we estimate

$$\tilde{\alpha}_{\text{comm}} = \sum_{\gamma_1, \gamma_2, \dots, \gamma_{p+1}} \left\| \left[H_{\gamma_{p+1}}, \dots, [H_{\gamma_2}, H_{\gamma_1}] \right] \right\| = \mathcal{O}(n).$$

- **We gave a product-formula algorithm with complexity $(nt)^{1+d/(\alpha-d)+o(1)}$** , outperforming the best previous approach based on Lieb-Robinson bounds.

Clustered Hamiltonian

Clustered Hamiltonian

$H = A + B = \sum_I H_I^{(1)} + \sum_I H_I^{(2)}$, where terms in A act on qubits within a single party and terms in B act between different parties.

- Appears naturally in classical fragmentation and Quantum Mechanics/Molecular Mechanics methods for large molecules.
- Group the terms within each party in A and simulate the resulting Hamiltonian using product formulas.¹³
- **Our new result implies a hybrid simulator with runtime $2^{\mathcal{O}\left(h_B^{\alpha(1)} t^{1+\alpha(1)} \text{cc}(g)/\epsilon^{\alpha(1)}\right)}$ with interaction strength h_B and contraction complexity $\text{cc}(g)$, improving the original result $2^{\mathcal{O}(h_B^2 t^2 \text{cc}(g)/\epsilon)}$.**

¹³[Peng, Harrow, Ozols, Wu 19]

Transverse field Ising model

Transverse field Ising model

$H = - \sum_{1 \leq u < v \leq n} j_{u,v} Z_u Z_v - \sum_{1 \leq u \leq n} h_u X_u$, where X_u, Z_u are Pauli operators acting on the u th qubit and $j_{u,v} \geq 0, h_u \geq 0$.

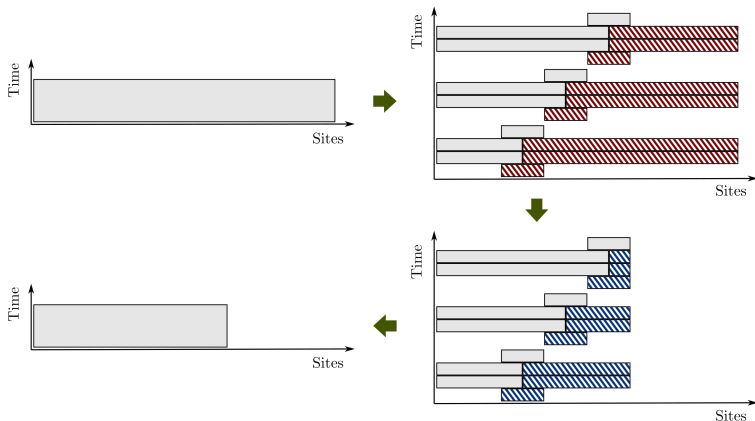
- The goal is to approximate the partition function $\text{Tr}(e^{-H})$ up to a multiplicative error.
- **We gave a Monte Carlo simulation of the transverse field Ising model with runtime $\tilde{O}(n^{45} j^{14} \epsilon^{-2} + n^{38} j^{21} \epsilon^{-9})$, tightening the previous result $\tilde{O}(n^{59} j^{21} \epsilon^{-9})$ of Bravyi.¹⁴**
- Similar improvement holds for the ferromagnetic quantum spin systems.¹⁵

¹⁴[Bravyi 15]

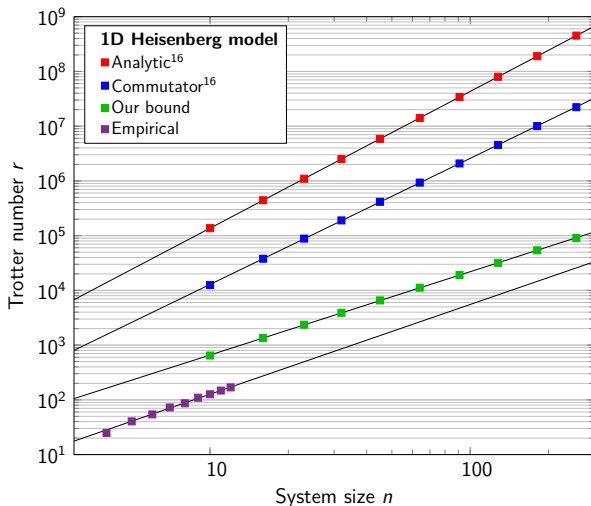
¹⁵[Bravyi, Gosset 17]

Simulating local observables

- We show that **local observables can be simulated with complexity independent of the system size** for power-law Hamiltonians, implying a Lieb-Robinson bound as a byproduct.



Tight prefactor



¹⁶[Childs, Maslov, Nam, Ross, Su 18]

A theory of Trotter error

Application	System	Best previous result	New result
Simulating quantum dynamics	Nearest-neighbor lattice	$(nt)^{1+\alpha(1)}$ (Conjecture), $\tilde{O}(nt)$ (Lieb-Robinson bound)	$(nt)^{1+\alpha(1)}$
	Electronic structure	$\tilde{O}(n^2t)$ (Interaction picture)	$n^{2+\alpha(1)}t^{1+\alpha(1)}$
	k -local Hamiltonians	$\tilde{O}(n^k \ H\ _1 t)$ (Qubitization)	$n^k \ H\ _1 \ H\ _1^{\alpha(1)} t^{1+\alpha(1)}$
	$1/x^\alpha$ ($\alpha > 2d$)	$\tilde{O}((nt)^{1+2d/(\alpha-d)})$ (Lieb-Robinson bound)	$(nt)^{1+d/(\alpha-d)+\alpha(1)}$
Simulating local observables	Clustered Hamiltonians	$2^{\mathcal{O}(h_B^2 t^2 cc(g)/\epsilon)}$	$2^{\mathcal{O}(h_B^{\alpha(1)} t^{1+\alpha(1)} cc(g)/\epsilon^{\alpha(1)})}$
Monte Carlo simulation	$1/x^\alpha$ ($\alpha > 2d$)	—	$t^{(1+d\frac{\alpha-d}{\alpha-2d})(1+\frac{d}{\alpha-d})+\alpha(1)}$
Monte Carlo simulation	Transverse field Ising model	$\tilde{O}(n^{59} j^{21} \epsilon^{-9})$	$\tilde{O}(n^{45} j^{14} \epsilon^{-2} + n^{38} j^{21} \epsilon^{-9})$
	Quantum ferromagnets	$\tilde{O}(n^{115}(1+\beta^{46})/\epsilon^{25})$	$\tilde{O}(n^{92}(1+\beta^{46})/\epsilon^{25})$

- Underpinning these improvements is a theory concerning the **types**, **order conditions**, and **representations** of Trotter error.

Error types

- Suppose that we use product formula $\mathcal{S}(t)$ to approximate the evolution of H for time $t \geq 0$.
- We consider the additive, exponentiated, and multiplicative type of Trotter error

$$\mathcal{S}(t) = e^{-itH} + \int_0^t d\tau e^{-i(t-\tau)H} \mathcal{S}(\tau) \mathcal{T}(\tau),$$

$$\mathcal{S}(t) = \exp_{\mathcal{T}} \left(-i \int_0^t d\tau (H + \mathcal{E}(\tau)) \right),$$

$$\mathcal{S}(t) = e^{-itH} \exp_{\mathcal{T}} \left(-i \int_0^t d\tau e^{i\tau H} \mathcal{E}(\tau) e^{-i\tau H} \right),$$

where $\mathcal{T}(\tau)$, $\mathcal{E}(\tau)$ consist of unitary conjugations of the form

$$e^{i\tau A_s} \dots e^{i\tau A_2} e^{i\tau A_1} B e^{-i\tau A_1} e^{-i\tau A_2} \dots e^{-i\tau A_s}.$$

Order conditions

- For a p th-order formula $\mathcal{S}_p(t)$, we have the following equivalent order conditions:
 - $\mathcal{S}_p(t) = e^{-itH} + O(t^{p+1})$;
 - $\mathcal{T}_p(\tau) = O(\tau^p)$; and
 - $\mathcal{E}_p(\tau) = O(\tau^p)$.
- Order conditions can be used to cancel lower-order terms in the Taylor expansion:

$$\mathcal{F}(\tau) = O(\tau^p) \iff \mathcal{F}(0) = \mathcal{F}'(0) = \dots = \mathcal{F}^{(p-1)}(0) = 0.$$

Error representations

- A unitary conjugation $e^{i\tau A_s} \dots e^{i\tau A_2} e^{i\tau A_1} B e^{-i\tau A_1} e^{-i\tau A_2} \dots e^{-i\tau A_s}$ has the expansion

$$C_0 + C_1\tau + \dots + C_{p-1}\tau^{p-1} + \mathcal{C}(\tau).$$

- Time-independent operators C_0, C_1, \dots, C_{p-1} can be canceled by order conditions.
- Operator-valued function $\mathcal{C}(\tau)$ is given by

$$\begin{aligned} \mathcal{C}(\tau) := & \sum_{j=1}^s \sum_{\substack{q_1 + \dots + q_j = p \\ q_j \neq 0}} e^{i\tau A_s} \dots e^{i\tau A_{j+1}} \\ & \cdot \int_0^\tau d\tau_2 e^{i\tau_2 A_j} \text{ad}_{iA_j}^{q_j} \dots \text{ad}_{iA_1}^{q_1}(B) e^{-i\tau_2 A_j} \cdot \frac{(\tau - \tau_2)^{q_j-1} \tau^{q_1 + \dots + q_{j-1}}}{(q_j - 1)! q_{j-1}! \dots q_1!} \\ & \cdot e^{-i\tau A_{j+1}} \dots e^{-i\tau A_s}, \end{aligned}$$

where $\text{ad}_{iA_1}(B) := [iA_1, B]$.

Outlook

- simulating time-dependent Hamiltonian $\sum_{\gamma=1}^{\Gamma} \mathcal{H}_{\gamma}(\tau)$;
- analysis of generalized product formulas (divide-and-conquer, randomized, LCU);
- improved circuit implementation for concrete systems;
- faster numerical computation of Trotter error bounds;
- different cost metric (e.g., sub-circuit model);
- simulating low-energy state;
- simulation in the presence of noise;
- other applications...



Analysis of the second-order formula

- For Hamiltonian $H = A+B$ and $t \geq 0$, the second-order formula $\mathcal{S}_2(t) = e^{-i\frac{t}{2}A}e^{-itB}e^{-i\frac{t}{2}A}$ can be represented as

$$\mathcal{S}_2(t) = e^{-itH} + \int_0^t d\tau_1 e^{-i(t-\tau_1)H} e^{-i\frac{\tau_1}{2}A} \mathcal{T}_2(\tau_1) e^{-\tau_1 B} e^{-i\frac{\tau_1}{2}A},$$

$$\mathcal{T}_2(\tau_1) = e^{-i\tau_1 B} \left(-i\frac{A}{2} \right) e^{i\tau_1 B} + i\frac{A}{2} + e^{i\frac{\tau_1}{2}A} (iB) e^{-i\frac{\tau_1}{2}A} - iB.$$

- We expand $\mathcal{T}_2(\tau_1)$ to second-order, obtaining

$$e^{-i\tau_1 B} \left(-i\frac{A}{2} \right) e^{i\tau_1 B} + i\frac{A}{2} = \left[-iB, -i\frac{A}{2} \right] \tau_1 + \int_0^{\tau_1} d\tau_2 \int_0^{\tau_2} d\tau_3 e^{-i\tau_3 B} \left[-iB, \left[-iB, -i\frac{A}{2} \right] \right] e^{i\tau_3 B},$$

$$e^{i\frac{\tau_1}{2}A} (iB) e^{-i\frac{\tau_1}{2}A} - iB = \left[i\frac{A}{2}, iB \right] \tau_1 + \int_0^{\tau_1} d\tau_2 \int_0^{\tau_2} d\tau_3 e^{i\frac{\tau_3}{2}A} \left[i\frac{A}{2}, \left[i\frac{A}{2}, iB \right] \right] e^{-i\frac{\tau_3}{2}A}.$$

Analysis of the second-order formula

- Altogether, we have the integral representation

$$\begin{aligned} & \mathcal{S}_2(t) - e^{-itH} \\ &= \int_0^t d\tau_1 \int_0^{\tau_1} d\tau_2 \int_0^{\tau_2} d\tau_3 e^{-i(t-\tau_1)H} e^{-i\frac{\tau_1}{2}A} \\ & \quad \cdot \left(e^{-i\tau_3 B} \left[-iB, \left[-iB, -i\frac{A}{2} \right] \right] e^{i\tau_3 B} + e^{i\frac{\tau_3}{2}A} \left[i\frac{A}{2}, \left[i\frac{A}{2}, iB \right] \right] e^{-i\frac{\tau_3}{2}A} \right) \\ & \quad \cdot e^{-\tau_1 B} e^{-i\frac{\tau_1}{2}A}, \end{aligned}$$

and the error bound

$$\left\| \mathcal{S}_2(t) - e^{-itH} \right\| \leq \frac{t^3}{12} \|[B, [B, A]]\| + \frac{t^3}{24} \|[A, [A, B]]\|.$$

- The general case follows by bootstrapping the above bound.

Trotter error with commutator scaling

- Consider a general p th-order product formula

$$\mathcal{S}_p(t) := \prod_{v=1}^{\Upsilon} \prod_{\gamma=1}^{\Gamma} e^{-ita_{(v,\gamma)} H_{\pi_v(\gamma)}} = e^{-itH} + O(t^{p+1}),$$

where $a_{(v,\gamma)}$ are real numbers with $|a_{(v,\gamma)}| \leq 1$.

Trotter error with commutator scaling

A p th-order formula $\mathcal{S}_p(t)$ can approximate the evolution of $H = \sum_{\gamma=1}^{\Gamma} H_{\gamma}$ for time $t \geq 0$ with Trotter error

$$\left\| \mathcal{S}_p(t) - e^{tH} \right\|, \left\| e^{-tH} \mathcal{S}_p(t) - I \right\| = O\left(\tilde{\alpha}_{\text{comm}} t^{p+1} e^{2t\Upsilon \sum_{\gamma=1}^{\Gamma} \|H_{\gamma}\|} \right),$$

where $\tilde{\alpha}_{\text{comm}} := \sum_{\gamma_1, \gamma_2, \dots, \gamma_{p+1}} \left\| \left[H_{\gamma_{p+1}}, \dots, [H_{\gamma_2}, H_{\gamma_1}] \right] \right\|$.