





Gauges, Loops, and Polynomials: for Partition Functions of Graphical Models

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Outline



Intro - I: Graphical Models

- Graphical Models: What?
- Graphical Models: Why?

2 Intro-II: Bethe Free Energy

- Variations: Exact \Rightarrow Approximate
- Poly- Representation for BFE & BP
- Belief Polytope. Forcing BP into Interior.
- Gauge T-, Graph R- & BP

- Gauge Transformation, Gauge Function & Loop Series
- Gauge Function ⇒ Partition Function via Graph/Algebraic R-
- Algebraic/Graph Reduction with BP- ... Exact & Approximate

Bi-Stability via BP & Gauges

- BP vs Exact
- Sequence of BP Low Bounds

Graphical Models: What? Graphical Models: Why?

What: Multi-Factor Graphical Model

- Multi-variate Probability Distribution: $p(\sigma) \doteq \frac{f(\sigma)}{Z}$
 - State vector: $\sigma = \{0,1\}^{|\mathcal{E}|}$
- Factorized according to the undirected graph, $(\mathcal{G}, \mathcal{E})$
 - components of σ reside on edges: $\forall \alpha \in \mathcal{E} : \sigma_{\alpha} = \{0, 1\}$
 - factors reside on nodes: $f(\sigma) \doteq \prod_{a \in \mathcal{V}} f_a(\sigma_a)$
 - $\forall a \in \mathcal{V}$: $\sigma_a \in \{0,1\}^{e_d(a)}$, e.g. $e_d(b) = (\gamma_d, \overline{\delta}_d, \overline{\theta}_d)$.
 - $(\mathcal{G}, \mathcal{E})$ includes self-loops & multi-edges



Graphical Models: What? Graphical Models: Why?

Why: GM for Optimization, Inference & Learning

Efficient Optimization & Inference

- Optimization = Maximal Likelihood
- Inference
 - Partition Function
 - Marginal Probabilities
 - Sampling

Efficient Learning = reconstruct GM from samples

Vuffray, Misra, Lokhov, MC (2016-) – not in this talk

Structure-specific OIL Applications

- Engineered Systems
 - Energy (Power, Natural Gas, Heating) Networks
- Physical Media
 - Fluid Mechanics

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Variations: Exact \Rightarrow Approximate Poly- Representation for BFE & BP Belief Polytope. Forcing BP into Interior.

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Exact Variational Principe for Partition Function

Gibbs (\leq 1902)-Kullblack-Leibler (1951)

• $\boldsymbol{b} = (0 \leq \boldsymbol{b}(\sigma) \leq 1 | \forall \sigma)$ - "belief" for each state

•
$$Z = \max_{\boldsymbol{b}} \prod_{\sigma} \left(\frac{\prod_{a} f_{a}(\sigma_{a})}{\boldsymbol{b}(\sigma)} \right)^{\boldsymbol{b}(\sigma)} \mid \forall \sigma : \boldsymbol{b}(\sigma) \ge 0$$

 $\sum_{\sigma} \boldsymbol{b}(\sigma) = 1$

- Convex optimization over exponentially large space
- Discrete \rightarrow continuous

Variations: Exact \Rightarrow Approximate Poly- Representation for BFE & BP Belief Polytope. Forcing BP into Interior.

Bethe- (BP-) Variational Principe for Partition Function

Yedidia-Freeman-Weiss (2005)

• Exact over tree-graphs

$$\begin{split} \boldsymbol{b}(\sigma) &\approx \frac{\prod_{a} \boldsymbol{b}_{a}(\sigma_{a})}{\prod_{\alpha \in \mathcal{E}} \boldsymbol{b}_{\alpha}(\sigma_{\alpha})} \text{ s.t.} \\ \forall \boldsymbol{a} \in \mathcal{V}, \ \forall \sigma_{\boldsymbol{a}} \in \boldsymbol{S}_{\boldsymbol{a}} : \ \boldsymbol{b}_{a}(\sigma_{a}) \doteq \sum_{\sigma \setminus \sigma_{a}} \boldsymbol{b}(\sigma) \\ \forall \alpha \in \mathcal{E}, \ \forall \sigma_{\alpha} = \{0,1\} : \ \boldsymbol{b}_{\alpha}(\sigma_{\alpha}) = \sum_{\sigma \setminus \sigma_{\alpha}} \boldsymbol{b}(\sigma) \end{split}$$

 Dynamic Programming ⇒ Belief Propagation Bethe (1935), Peierls (1935), Gallager (1961), Pearl (1988)

• Ansatz in general. Substitute into Gibbs-Kubblack-Leibler \Rightarrow

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Bethe- (BP-) Variational Principe for Partition Function

Yedidia-Freeman-Weiss (2005)

$$\min_{b}(-\log \mathcal{Z}(b)) \text{ s.t. } \begin{cases} \forall a \in \mathcal{V} \quad \forall \alpha \in e_{d}(a), \ \forall \tau \in \{0,1\}: \quad b_{\alpha}(\tau) = \sum_{s \in S_{a}}^{s_{\alpha} = \tau} b_{a}(s) \\ \forall s \in S_{a}: \qquad \qquad b_{a}(s) \geq 0 \\ \sum_{s \in S_{a}} b_{a}(s) = 1 \end{cases}$$

where $\mathcal{Z}(b)$ is the Belief Propagation Partition Function (BP-PF)

$$\mathcal{Z}(b)\doteq \left(\prod_{a\in\mathcal{V}}\prod_{\sigma_a\in\mathcal{S}_a}\left(rac{f_a(\sigma_a)}{b_a(\sigma_a)}
ight)^{b_a(\sigma_a)}
ight)\left(\prod_{lpha\in\mathcal{E}}\prod_{\sigma_lpha\in\{0,1\}}\left(b_lpha(\sigma_lpha)
ight)^{b_lpha(\sigma_lpha)}
ight)^{b_a(\sigma_a)}
ight)$$

of the vector b of marginal beliefs associated with edges and nodes

$$b = (b_{\alpha}(\sigma_{\alpha})| \alpha \in \mathcal{E}, \ \sigma_{\alpha} \in \{0,1\}) \oplus (b_{a}(\sigma_{a})| a \in \mathcal{V}, \ \sigma_{a} \in S_{a}).$$

Variations: Exact \Rightarrow Approximate Poly- Representation for BFE & BP Belief Polytope. Forcing BP into Interior.

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Poly- (Dual-) Representation (for Bethe Appr. of Z)

Anari & Oveis Gharan 2017, Straczak & Vishnoi 2017

$$\begin{split} \sup_{\beta \in [0;1]^{\mathcal{E}}} \min_{x \in \mathbb{R}_{+}^{\mathcal{E}_{d}}} \mathcal{L}(\beta, x) \\ \mathcal{L}(\beta, x) \doteq \left(\prod_{\alpha \in \mathcal{E}} \beta_{\alpha}^{\beta_{\alpha}} (1 - \beta_{\alpha})^{1 - \beta_{\alpha}} \right) \prod_{a \in \mathcal{V}} \frac{h_{a}(x_{a})}{\prod_{\alpha \in \mathcal{E}_{d}(a)} x_{\alpha}^{\beta_{\alpha}}} \\ \forall a \in \mathcal{V} : \quad h_{a}(x_{a}) \doteq \sum_{s \in S_{a}} f_{a}(s) \prod_{\alpha \in e_{d}(a)} x_{\alpha}^{s_{\alpha}} \end{split}$$

• Generalizes Gurvits (2011 – permanent)

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Graphical Model Examples: Hard- & Soft-

Hard Example ($\forall a : \exists \sigma_a \text{ s.t. } f_a(\sigma_a) = 0$): Perfect Matching

•
$$\forall a: f_a(\sigma_a) = \begin{cases} \prod_{\alpha \in e_d(a)} (\mu_\alpha)^{\sigma_\alpha}, & \sum_{\alpha \in e_d(a)} \sigma_\alpha = 1 \\ 0, & \text{otherwise} \end{cases}$$

- Bethe (Free Energy) optimization is convex (Vontobel 2010)
- BP solution may be on the BP-polytope boundary (at low "temperature") (Watanabe, MC 2009)
- In the case of bi-partite graph Bethe appr. gives low bound for Z (Gurvits 2011)

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Graphical Model Examples: Hard- & Soft-

Soft Example (= $\forall a : f_a(\sigma_a) > 0$): "soft" Perfect Matching

•
$$\forall a: f_a(\sigma_a) = \begin{cases} \prod_{\alpha \in e_d(a)} (\mu_\alpha)^{\sigma_\alpha}, \sum_{\alpha \in e_d(a)} \sigma_\alpha = 1 \\ \epsilon > 0, & \text{otherwise} \end{cases}$$

For any Soft GM

• <u>Theorem</u>: Solution of Bethe optimization is always in the interior of BP polytope (MC, VC, YM 2019)

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Graphical Model Examples: Hard- & Soft-

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For any Soft GM

 <u>Theorem</u>: Solution of Bethe optimization is always in the interior of BP polytope (MC, VC, YM 2019)

will assume small- ϵ regularization for the remainder of the talk

• BP (= solution of Bethe optimization) is in the interior

Gauge Transformation, Gauge Function & Loop Series Gauge Function \Rightarrow Partition Function via Graph/Algebraic R-Algebraic/Graph Reduction with BP- ... Exact & Approximate

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Gauge Transformation

[MC, Chernyak 2006]

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Gauge-T = multi-linear transformation of factors:

•
$$\forall a, \forall \sigma_a : f_a(\sigma_a) \to \sum_{\varsigma_a} f_a(\varsigma_a) \prod_{\alpha_d \in e_d(a)} G_{\alpha_d}(\sigma_\alpha, \varsigma_{\alpha_d})$$

which keeps Z invariant

•
$$\forall G: Z = \sum_{\sigma} \prod_{a \in \mathcal{V}} f_a(\sigma_a) = \sum_{\sigma} z(\sigma|G)$$

•
$$z(\sigma|G) = \prod_{a} f_{a}(\varsigma_{a}) \prod_{\alpha_{d} \in e_{d}(a)} G_{\alpha_{d}}(\sigma_{\alpha}, \varsigma_{\alpha_{d}})$$

•
$$\varsigma_{a} \doteq (\varsigma_{\alpha_{d}} = 0, 1 | \alpha_{d} \in e_{d}(a))$$

related approaches

- Reparametrization (Wainwright, Jaakola, Willsky 2003)
- Hollographic transformation/algorithm (Valliant 2004)

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$$\forall G: \quad Z = \sum_{\sigma} \prod_{a \in \mathcal{V}} f_a(\sigma_a) = \sum_{\sigma} z(\sigma | G)$$

• $z(\sigma | G) \doteq \prod_a f_a(\varsigma_a) \prod_{\alpha d \in e_d(a)} G_{\alpha_d}(\sigma_\alpha, \varsigma_{\alpha_d})$

•
$$\varsigma_{a} \doteq (\varsigma_{\alpha_{d}} = 0, 1 | \alpha_{d} \in e_{d}(a))$$

Orthogonality of GT

• **G** are (2 × 2) matrices two per age

•
$$\forall \alpha \in \mathcal{E}, \quad \mathbf{G}_{\alpha_d}^{\mathcal{T}} * \mathbf{G}_{\bar{\alpha}_d} = 1$$

guarantees invariance of Z



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 - $\forall G: Z = \sum_{\sigma} \prod_{a \in \mathcal{V}} f_a(\sigma_a) = \sum_{\sigma} z(\sigma|G)$
 - $z(\sigma|G) \doteq \prod_{a} f_{a}(\varsigma_{a}) \prod_{\alpha_{d} \in e_{d}(a)} G_{\alpha_{d}}(\sigma_{\alpha}, \varsigma_{\alpha_{d}})$

•
$$\varsigma_{a} \doteq (\varsigma_{\alpha_{d}} = 0, 1 | \alpha_{d} \in e_{d}(a))$$

Polynomial Representation of Gauges

•
$$\forall \alpha_d \in \mathcal{E}_d : \quad G_{\alpha_d} =$$

$$\frac{1}{(x_{\alpha_d} \times \bar{\alpha}_d)^{1/4} \sqrt{1 + x_{\alpha_d} \times \bar{\alpha}_d}} \begin{pmatrix} \sqrt{x_{\bar{\alpha}_d}} & x_{\alpha_d} \sqrt{x_{\bar{\alpha}_d}} \\ -x_{\bar{\alpha}_d} \sqrt{x_{\alpha_d}} & \sqrt{x_{\alpha_d}} \end{pmatrix}$$
• $x \doteq (x_{\alpha_d} > 0 | \alpha_d \in \mathcal{E}_d)$ is positive component-wise

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$$\forall G : Z = \sum_{\sigma} \prod_{a \in \mathcal{V}} f_a(\sigma_a) = \sum_{\sigma} z(\sigma | G)$$

• $z(\sigma | G) \doteq \prod_a f_a(\varsigma_a) \prod_{\alpha \neq e_d(a)} G_{\alpha_d}(\sigma_\alpha, \varsigma_{\alpha_d})$

•
$$\varsigma_a \doteq (\varsigma_{\alpha_d} = 0, 1 | \alpha_d \in e_d(a))$$

Theorem: Any extremum of BFE (a BP solution) is a stationary point of z(x). Minimum of BFE is achieved at the extremum maximizing z(x).

$$\begin{aligned} & \forall a \in \mathcal{V}, \quad \forall \alpha \in \mathcal{E}_{\mathrm{d}}(a) : \left. \partial_{x_{\alpha}} z(x) \right|_{x=x^{(\mathrm{bp})}} = 0 \\ & z(x) \doteq \frac{\prod\limits_{a \in \mathcal{V}} h_{a}(x_{a})}{\prod\limits_{\alpha \in \mathcal{E}} (1+x_{\alpha_{d}} x_{\tilde{\alpha}_{d}})} = \frac{h(x)}{\prod\limits_{\alpha \in \mathcal{E}} (1+x_{\alpha_{d}} x_{\tilde{\alpha}_{d}})} \end{aligned}$$

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•
$$\varsigma_{a} \doteq (\varsigma_{\alpha_{d}} = 0, 1 | \alpha_{d} \in e_{d}(a))$$

Loop Series/Calculus

$$\begin{split} & Z = \sum_{\sigma} z(\sigma | x^{(bp)}), \quad \forall \sigma : \quad z(\sigma | x^{(bp)}) = Z^{(bp)} \frac{\prod_{a} \mu_{a}^{(bp)}}{\prod_{\alpha} \beta_{\alpha}^{(bp)} (1 - \beta_{\alpha}^{(bp)})} \\ & \forall a : \quad \mu_{a} \doteq \frac{\sum_{\varsigma_{a}} f_{a}(\varsigma_{a}) \prod_{\alpha} \left((x_{\alpha}^{(bp)})^{\varsigma_{\alpha}} (\varsigma_{\alpha} - \mu_{\alpha})^{\sigma_{\alpha}} \right)}{\sum_{\varsigma_{a}} f_{a}(\varsigma_{a}) \prod_{\alpha \in e} (\sigma)_{(a)} (x_{\alpha}^{(bp)})^{\varsigma_{\alpha}}} \end{split}$$



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• $z(\sigma|G) \doteq \prod_{a \in \mathcal{V}} f_a(\sigma_a) = \sum_{\sigma} z(\sigma|G)$

$$\int 2(O | \mathbf{G}) = \prod_{a} I_{a}(\varsigma_{a}) \prod_{\alpha_{d} \in e_{d}(a)} \mathbf{G}_{\alpha_{d}}(O_{\alpha}, \varsigma_{\alpha_{d}})$$

•
$$\varsigma_a \doteq (\varsigma_{\alpha_d} = 0, 1 | \alpha_d \in e_d(a))$$

Further (finite graph) uses of "Beyond BP" = Loop Series & Gauge T-

- Minimum Weight Perfect Matching (optimization) Blossom BP [Sungsoo Ahn, Sejun Park, Jinwoo Shin & MC NIPS 2015]
- LS maps GM to another GM (possibly with negative weights). The new GM can be sampled with MCMC. [Sungsoo Ahn, Jinwoo Shin, MC NIPS 2016]
- BP can be mixed with Mean Field Gauging Variational Inference [Sungsoo Ahn, Jinwoo Shin & MC NIPS 2017]

Gauge Transformation, Gauge Function & Loop Series Gauge Function \Rightarrow Partition Function via Graph/Algebraic R-Algebraic/Graph Reduction with BP- ... Exact & Approximate

Partition Function via Graph-R = Elimination of Edges



- Fix edge-elimination order
- Sum over edge variables sequentially
- Get a sequence of GMs
 - node-degree & complexity grows (exponentially)

Proposition:

Partition Functions in Graph-R sequence are identical/invariant

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From Gauge-F to Partition Function

Gauge Function

$$z(x) \doteq \frac{\prod\limits_{a \in \mathcal{V}} h_a(x_a)}{\prod\limits_{\alpha \in \mathcal{E}} (1 + x_{\alpha_d} x_{\bar{\alpha}_d})}$$

<u>Theorem</u>: The Algebraic (Differentiate & Marginalize) -T applied to the Gauge-F results in Z

$$\begin{split} m &= 0: \quad x^{(0)} \doteq x, \quad \mathcal{Z}^{(0)}(x) \doteq \mathbf{z}(x) \\ m &= 1, \cdots, |\mathcal{E}|: \quad x^{(m)} \doteq x^{(m-1)} \setminus \{x_{\alpha_d^{(m)}}, x_{\bar{\alpha}_d^{(m)}}\} \\ \mathcal{Z}^{(m)}(x^{(m)}) \doteq \left(1 + \partial_{x_{\alpha_d^{(m)}}} \partial_{x_{\bar{\alpha}_d^{(m)}}}\right) \left(\left(1 + x_{\alpha_d^{(m)}} x_{\bar{\alpha}_d^{(m)}}\right) \mathcal{Z}^{(m-1)} \right) \Big|_{x_{\alpha_d^{(m)}} = x_{\bar{\alpha}_d^{(m)}} = 0} \\ m &= |\mathcal{E}|: \quad x^{(|\mathcal{E}|)} = \emptyset, \quad \mathcal{Z}^{(|\mathcal{E}|)} = \mathbf{Z} \end{split}$$

Gauge Transformation, Gauge Function & Loop Series Gauge Function \Rightarrow Partition Function via Graph/Algebraic R-Algebraic/Graph Reduction with BP- ... Exact & Approximate

Graph-R is equivalent to Algebraic-R



Theorem:Algebraic (diff.+marg.) -R = Graph (edge elimination)-T

$$\begin{aligned} \mathcal{Z}^{(m)}(x^{(m)}) &= \frac{h^{(m)}(x^{(m)})}{\prod_{\alpha \in \mathcal{E}^{(m)}} (1 + x_{\alpha_d} x_{\bar{\alpha}_d})} \\ h^{(m)}(x^{(m)}) &\doteq \prod_{a \in \mathcal{V}^{(m)}} \left(\sum_{\varsigma_a} f_a^{(m)}(\varsigma_a) \prod_{\alpha_d \in \mathfrak{e}_d(a)} x_{\alpha_d}^{\varsigma_{\alpha_d}} \right) \end{aligned}$$

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Edge Elimination with Belief Propagation

Theorem: Any extremum of BFE (a BP solution) is a stationary point of z(x). Minimum of BFE is achieved at the extremum maximizing z(x).

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Extremum of BFE:

Gauge Function:

 $\begin{aligned} \forall \mathbf{a} \in \mathcal{V}, \quad \forall \alpha \in \mathcal{E}_{\mathrm{d}}(\mathbf{a}) : \left. \partial_{\mathbf{x}_{\alpha}} z(\mathbf{x}) \right|_{\mathbf{x} = \mathbf{x}^{(\mathrm{bp})}} &= \mathbf{0} \\ z(\mathbf{x}) \doteq \frac{\prod\limits_{a \in \mathcal{V}} h_{a}(\mathbf{x}_{a})}{\prod\limits_{\alpha \in \mathcal{E}} \left(1 + \mathbf{x}_{\alpha_{d}} \mathbf{x}_{\tilde{\alpha}_{d}}\right)} &= \frac{h(\mathbf{x})}{\prod\limits_{\alpha \in \mathcal{E}} \left(1 + \mathbf{x}_{\alpha_{d}} \mathbf{x}_{\tilde{\alpha}_{d}}\right)} \end{aligned}$

Let us "eliminate" a single edge

$$\begin{split} h(x) &\doteq h^{(0,0)} + h^{(1,0)} x_{\alpha_d} + h^{(0,1)} x_{\bar{\alpha}_d} + h^{(1,1)} x_{\alpha_d} x_{\bar{\alpha}} \\ \partial_{x_{\alpha_d}} \partial_{x_{\bar{\alpha}_d}} \frac{h(x)}{1 + x_{\alpha_d} x_{\bar{\alpha}_d}} = 0 \Big|_{x^{(\alpha - bp)}} \quad \Rightarrow \end{split}$$

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Let us "eliminate" a single edge

$$\begin{split} h(x) &\doteq h^{(0,0)} + h^{(1,0)} x_{\alpha_d} + h^{(0,1)} x_{\bar{\alpha}_d} + h^{(1,1)} x_{\alpha_d} x_{\bar{\alpha}_d} \\ \partial_{x_{\alpha_d}} \partial_{x_{\bar{\alpha}_d}} \frac{h(x)}{1 + x_{\alpha_d} x_{\bar{\alpha}_d}} &= 0 \Big|_{x^{(\alpha - bp)}} \Rightarrow \\ 2x^{(\alpha - bp)}_{\alpha_d} h^{(1,0)} &= 2x^{(\alpha - bp)}_{\bar{\alpha}_d} h^{(0,1)} = h^{(1,1)} - h^{(0,0)} + \sqrt{\left(h^{(1,1)} - h^{(0,0)}\right)^2 + 4h^{(0,1)} h^{(1,0)}} \end{split}$$

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$$\begin{aligned} \forall a \in \mathcal{V}, \quad \forall \alpha \in \mathcal{E}_{\mathrm{d}}(a) : \left. \partial_{x_{\alpha}} z(x) \right|_{x=x^{\mathrm{(bp)}}} &= 0 \\ z(x) \doteq \frac{\prod\limits_{a \in \mathcal{V}} h_a(x_a)}{\prod\limits_{\alpha \in \mathcal{E}} (1+x_{\alpha_d} x_{\tilde{\alpha}_d})} &= \frac{h(x)}{\prod\limits_{\alpha \in \mathcal{E}} (1+x_{\alpha_d} x_{\tilde{\alpha}_d})} \end{aligned}$$

Let us "eliminate" a single edge

$$\begin{split} h(x) \doteq h^{(0,0)} + h^{(1,0)} x_{\alpha_d} + h^{(0,1)} x_{\bar{\alpha}_d} + h^{(1,1)} x_{\alpha_d} x_{\bar{\alpha}_d} \\ \partial_{x_{\alpha_d}} \partial_{x_{\bar{\alpha}_d}} \frac{h(x)}{1 + x_{\alpha_d} x_{\bar{\alpha}_d}} = 0 \Big|_{x^{(\alpha - bp)}} \Rightarrow \\ & 2x^{(\alpha - bp)}_{\alpha_d} h^{(1,0)} = 2x^{(\alpha - bp)}_{\bar{\alpha}_d} h^{(0,1)} = h^{(1,1)} - h^{(0,0)} + \sqrt{\left(h^{(1,1)} - h^{(0,0)}\right)^2 + 4h^{(0,1)}h^{(1,0)}} \\ & \underline{\text{If } \alpha \text{ is a normal edge}} \text{ BP Graph Reduction (edge contraction) is exact} \\ & x^{(\alpha - bp)}_{\alpha_d} = \frac{h^{(1,1)}}{h^{(1,0)}}, \quad x^{(\alpha - bp)}_{\bar{\alpha}_d} = \frac{h^{(1,1)}}{h^{(0,1)}}, \quad h^{(\alpha - bp)} = h^{(1,1)} + h^{(0,0)} \end{split}$$

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Gauge Transformation, Gauge Function & Loop Series Gauge Function \Rightarrow Partition Function via Graph/Algebraic R-Algebraic/Graph Reduction with BP- ... Exact & Approximate

Graph Reduction with BP to Bouquet (exact)



Introduce BP at each step in a Graph-R sequence

• Theorem: BP Graph-R of a normal edge is exact

Definition: Bouquet-graph: contract all normal edges

• **Corollary**: sequence of BP Graph-R leading to bouquet-graph is exact

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Graph Reduction with BP



- BP-sequence (in general) is not monotonic.
- Complexity increases with contraction.
- Can be stopped/truncated at any step
 - \Rightarrow an approximate elimination

Gauge Transformation, Gauge Function & Loop Series Gauge Function ⇒ Partition Function via Graph/Algebraic R-Algebraic/Graph Reduction with BP- ... Exact & Approximate

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No assumptions about GM so far – fully general



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BP vs Exact Sequence of BP Low Bounds

Outline



1 Intro - I: Graphical Models

- Graphical Models: What?
- Graphical Models: Why?

2 Intro-II: Bethe Free Energy

- Variations: Exact \Rightarrow Approximate
- Poly- Representation for BFE & BP
- Belief Polytope. Forcing BP into Interior.
- 3) Gauge T-, Graph R- & BF

- Gauge Transformation, Gauge Function & Loop Series
- Gauge Function ⇒ Partition
 Function via Graph/Algebraic R-
- Algebraic/Graph Reduction with BP- ... Exact & Approximate

Bi-Stability via BP & Gauges

- BP vs Exact
- Sequence of BP Low Bounds

BP vs Exact Sequence of BP Low Bounds

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BP vs Exact for an Edge-elimination step

$$\begin{aligned} \forall a \in \mathcal{V}, \quad \forall \alpha \in \mathcal{E}_{d}(a) : \left. \partial_{x_{\alpha}} z(x) \right|_{x=x^{(bp)}} &= 0\\ z(x) \doteq \frac{\prod_{a \in \mathcal{V}} h_{a}(x_{a})}{\prod_{\alpha \in \mathcal{E}} (1 + x_{\alpha_{d}} x_{\bar{\alpha}_{d}})} &= \frac{h(x)}{\prod_{\alpha \in \mathcal{E}} (1 + x_{\alpha_{d}} x_{\bar{\alpha}_{d}})}\\ h(x) \doteq h^{(0,0)} + h^{(1,0)} x_{\alpha_{d}} + h^{(0,1)} x_{\bar{\alpha}_{d}} + h^{(1,1)} x_{\alpha_{d}} x_{\bar{\alpha}_{d}} \end{aligned}$$

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BP vs Exact Sequence of BP Low Bounds

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BP vs Exact for an Edge-elimination step

$$\begin{aligned} \forall a \in \mathcal{V}, \quad \forall \alpha \in \mathcal{E}_{\mathrm{d}}(a) : \partial_{x_{\alpha}} z(x)|_{x=x^{(\mathrm{bp})}} = 0 \\ z(x) \doteq \frac{\prod\limits_{a \in \mathcal{V}} h_a(x_a)}{\prod\limits_{\alpha \in \mathcal{E}} (1+x_{\alpha_d} x_{\bar{\alpha}_d})} = \frac{h(x)}{\prod\limits_{\alpha \in \mathcal{E}} (1+x_{\alpha_d} x_{\bar{\alpha}_d})} \\ y(x) \doteq h^{(0,0)} + h^{(1,0)} x_{\alpha_d} + h^{(0,1)} x_{\bar{\alpha}_d} + h^{(1,1)} x_{\alpha_d} x_{\bar{\alpha}} \end{aligned}$$

Lemma: BP Reduction vs Exact Reduction (Differential Version)

$$\begin{aligned} \forall x^{(1)} &= x \setminus \{x_{\alpha_d}, x_{\bar{\alpha}_d}\} : \left. \frac{h(x)}{1 + x_{\alpha_d} \times \bar{\alpha}_d} \right| & x_{\alpha_d} = x_{\alpha_d}^{(\alpha - bp)} & \leq h^{(1,1)}(x^{(1)}) + h^{(0,0)}(x^{(1)}) \\ & x_{\bar{\alpha}_d} = x_{\bar{\alpha}_d}^{(\alpha - bp)} \\ & \underline{holds if} & h^{(0,1)}(x^{(1)})h^{(1,0)}(x^{(1)}) \leq h^{(0,0)}(x^{(1)})h^{(1,1)}(x^{(1)}) \end{aligned}$$

BP vs Exact Sequence of BP Low Bounds

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BP vs Exact for an Edge-elimination step

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<u>Lemma:</u> BP Reduction vs Exact Reduction (Variational Version) [Anari & Oveis Gharan 2017] & [Straczak & Vishnoi 2017]

 $\forall x^{(1)}: \quad h^{(0,1)}(x^{(1)})h^{(1,0)}(x^{(1)}) \leq h^{(0,0)}(x^{(1)})h^{(1,1)}(x^{(1)})$ guarantees that

$$\forall \beta_{\alpha} \in [0;1], \quad (\beta_{\alpha})^{\beta_{\alpha}} (1-\beta_{\alpha})^{1-\beta_{\alpha}} \inf_{x_{\alpha_{d}}, x_{\bar{\alpha}_{d}}} \frac{h(x)}{(x_{\alpha_{d}} x_{\bar{\alpha}_{d}})^{\beta_{\alpha}}} \leq h^{(1,1)}(x^{(1)}) + h^{(0,0)}(x^{(1)})$$

BP vs Exact Sequence of BP Low Bounds

When $Z^{(bp)}$ is a low bound?

Definition: [Real-Stability (RS) and Bi-Stability (BS)]

- A nonzero polynomial, $g(x) \in \mathbb{R}[x_1, \dots, x_N]$, with real coefficients is RS if none of its roots $z = (z_1, \dots, z_N) \in \mathbb{C}^N$ satisfies: $\operatorname{Im}(z_i) > 0$ for every $i = 1, \dots, N$.
- A polynomial $h(x_{\alpha^{(1)}}, x_{\bar{\alpha}_d}^{(1)}; x_{\alpha^{(2)}}, x_{\bar{\alpha}_d}^{(2)} \cdots)$ is BS if $h(x_{\alpha^{(1)}}, -x_{\bar{\alpha}_d}^{(1)}; x_{\alpha^{(2)}}, -x_{\bar{\alpha}_d}^{(2)} \cdots)$ is Real Stable.

Theorem: Monotonicity of BP-estimations (for contraction sequence)

If $\forall a \in \mathcal{V}, h_a(x_a)$ is bi-stable then

- each polynomial in the contraction (=differentiate+marginalize) sequence, $h^{(m)}(x^{(m)})$, is stable
- Value of BP estimation for Z does not decrease with elimination: $Z^{(bp)} = Z^{(0;bp)} < Z^{(1;bp)} < \dots < Z^{(|\mathcal{E}|;bp)} = Z$

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BP vs Exact Sequence of BP Low Bounds

When $Z^{(bp)}$ is a low bound?

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$Z^{(bp)} \leq Z$

Gurvits (2011) - permanents
 Straczak & Vishnoi 2017 - bi-partite stable
 Anari & Oveis Gharan 2017 - bi-stable

BP vs Exact Sequence of BP Low Bounds

When $Z^{(bp)}$ is a low bound?

$Z^{(bp)} \leq Z$

 Attractive (ferro) Ising model via Loop Series, with restrictions (all loop terms are positive) Sudderth, Wainwright, Willsky 2007

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• via Graph Covers (Vontobel) with no restrictions Ruozi 2012

another approach - not related to the poly- considerations

BP vs Exact Sequence of BP Low Bounds

Summary



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BP vs Exact Sequence of BP Low Bounds

Path Forward

- Mixing the "differentiate and marginalize" approach with the Loop Calculus/Series
- Generalization to higher alphabets (homogeneous polynomials)
- **Tighter** Approximations (mix of gauge/BP- and graph-transformations)
 - theory & heuristics
- From global to local stability broader class of "tractable" polynomials
- Sequential Elimination, e.g. extending
 - Gauged Mini-Bucket Elimination for Approximate Inference [Sungsoo Ahn, Jinwoo Shin, Adrian Weller & MC 2018]
 - Bucket renormalization for approximate inference [Sungsoo Ahn, Jinwoo Shin, Adrian Weller & MC 2018]

BP vs Exact Sequence of BP Low Bounds



Thank You !

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