

Extremal Mechanisms in Differential Privacy



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Information Theory and Differential Privacy



- Communication -- **small** error probability
- Privacy -- **large** error probability

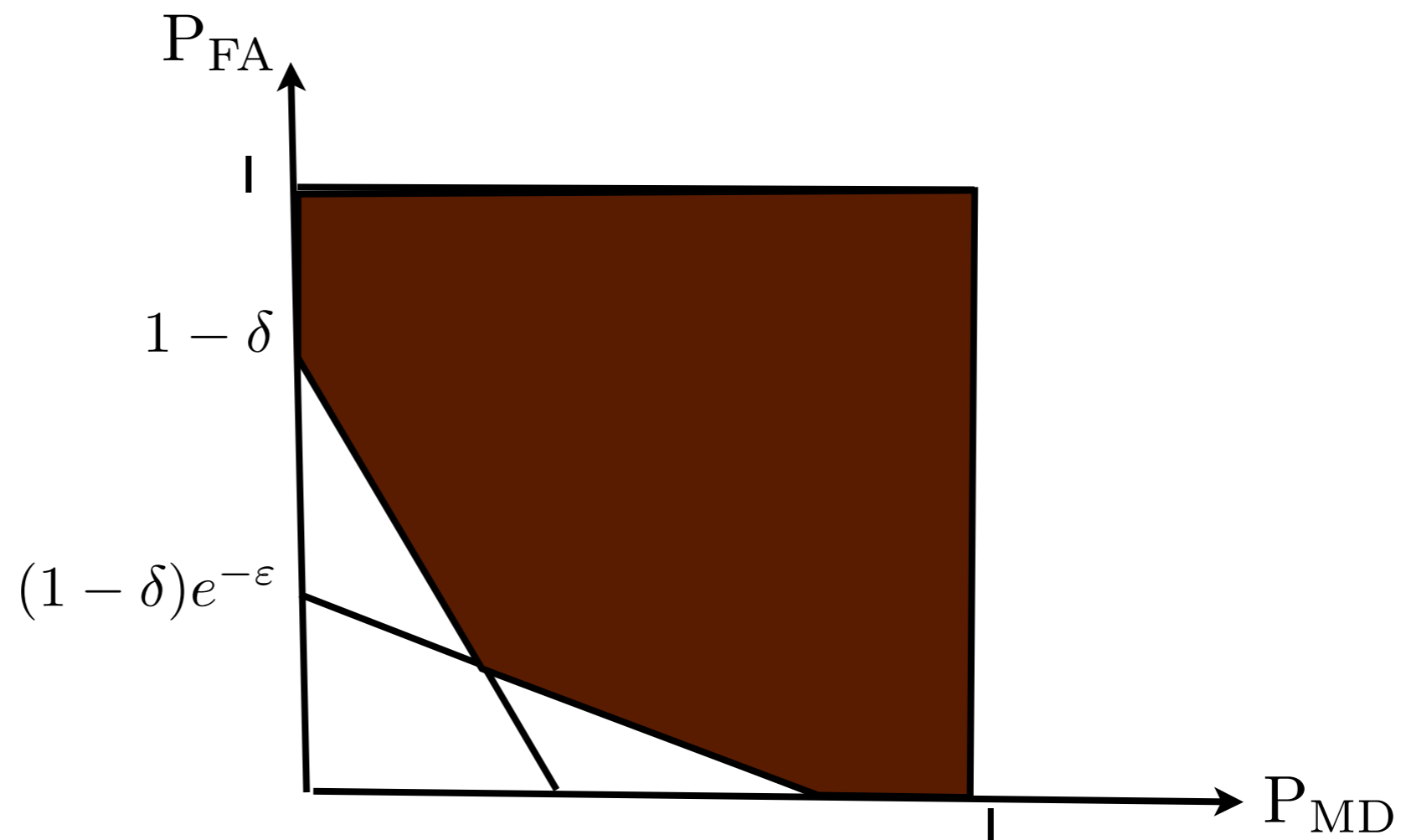
Information Theory and Differential Privacy



- Communication -- multi hypothesis testing
- Privacy -- binary hypothesis testing

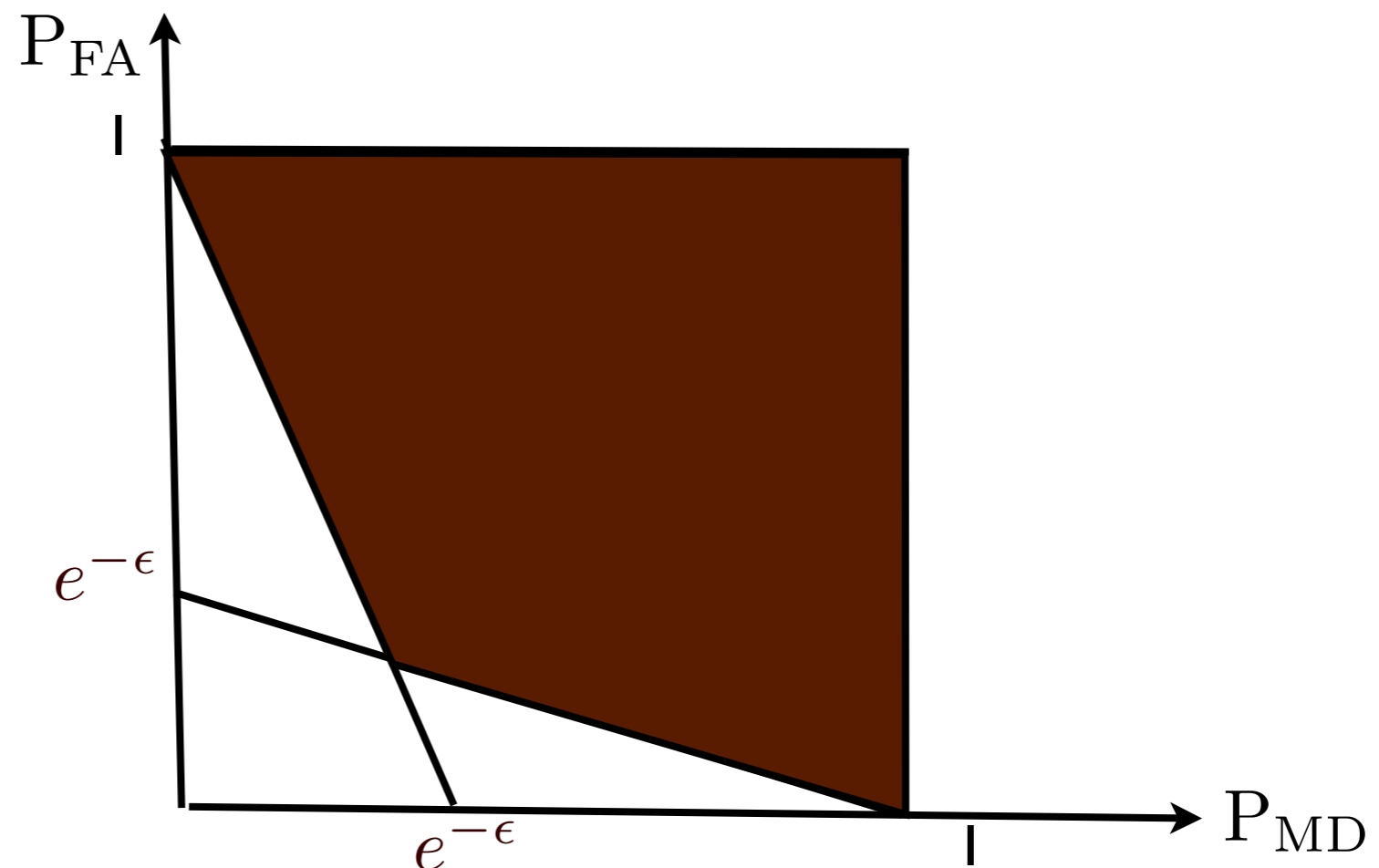
Binary Inference Errors

- **Two error** types
 - **False Alarm** and **Missed Detection**
- **Privacy: guarantee enough error**



Differential Privacy

- A specific way of enforcing inference errors
 - WZ11
- Original formulation involves likelihood ratios
 - DKMNS05
- ϵ controls privacy level



•

Differential Privacy

- For competing hypotheses D_1 and D_2

$$e^{-\epsilon} \leq \frac{\Pr(K(D_1) \in S)}{\Pr(K(D_2) \in S)} \leq e^{\epsilon}$$

- Equivalently:

$$P_{\text{MD}} + e^{-\epsilon} P_{\text{FA}} \geq e^{-\epsilon}$$

$$P_{\text{FA}} + e^{-\epsilon} P_{\text{MD}} \geq e^{-\epsilon}$$

- Likelihood ratios in a bounded interval
- ϵ small is high privacy
- ϵ large is low privacy

Information Theory is Mature



- Shannon, 1948
 - A mathematical theory of communication
- Success
 - extremal limits
 - capacity, single-letter expressions
 - fundamental benchmarks
 - practical schemes
 - operational interpretation
 - data processing inequalities

This Talk

- Similar program for differential privacy
 - extremal mechanisms
 - fundamental limits
 - operational interpretation
- Results
 - Staircase mechanism
 - universally optimal noise adding mechanism
 - Optimal Composition theorems
 - Abstract Staircase mechanism
 - dominates every other privacy mechanism

State of the Art

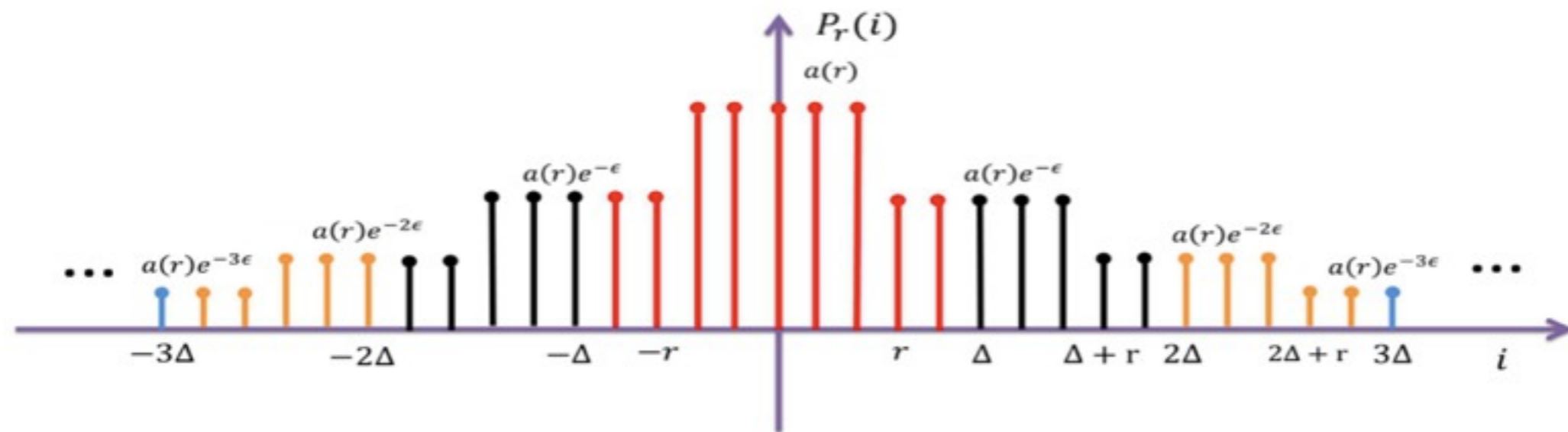
- Noise adding mechanisms
- Real valued query
 - Laplacian noise
 - regular differential privacy
 - Gaussian noise
 - approximate differential privacy
- No exact optimality results

State of the Art

- Integer valued query
- Count queries (sensitivity is one)
- Geometric noise added
 - universal optimality in Bayesian cost minimization framework [GRS09]
 - no natural generalization
 - larger sensitivity [GS10]
- No operational interpretation
 - Hint: Log Likelihood ratio $\in \{-\varepsilon, +\varepsilon\}$

Staircase Mechanism

- **Universally optimal noise** adding mechanism
 - worst case setting
 - generalization of GRS09 ($\Delta = 1$)



- no operational interpretation
 - Log Likelihood ratio $\in \{-\varepsilon, 0, +\varepsilon\}$

Example Cost Functions

- Privacy mechanism involves adding noise

$$K(D) = q(D) + X$$

- **amplitude of noise** $E[|X|]$ $L(x) = |x|$

- **variance of noise** $E[X^2]$ $L(x) = x^2$

- In general any cost function

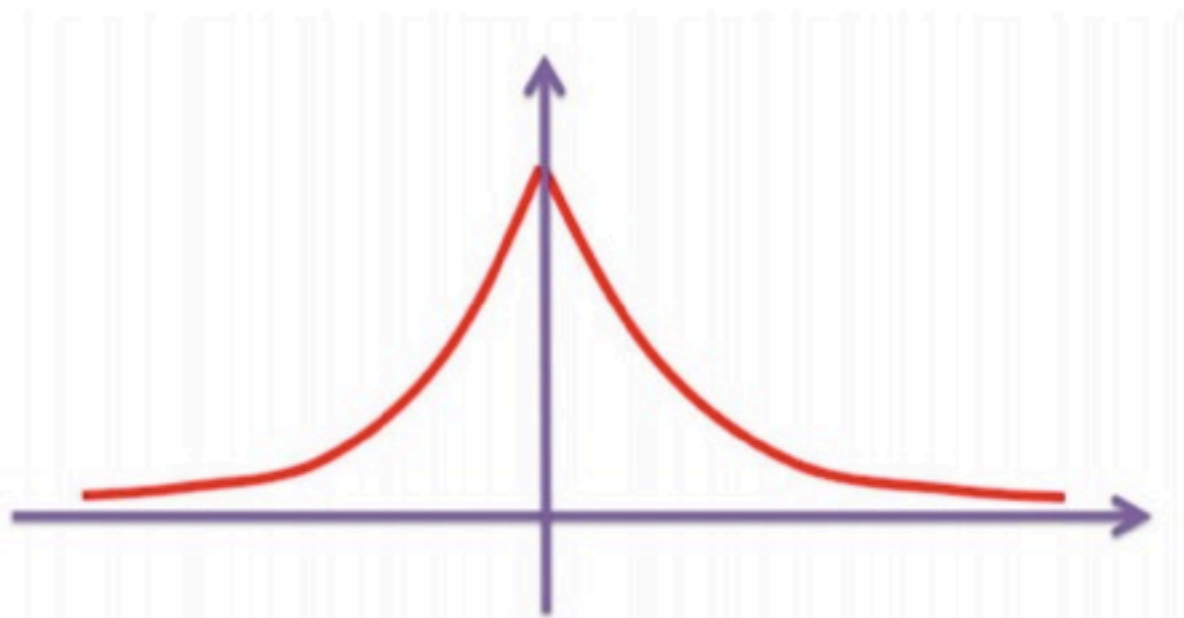
- **monotonically increasing**

- **symmetric around origin**

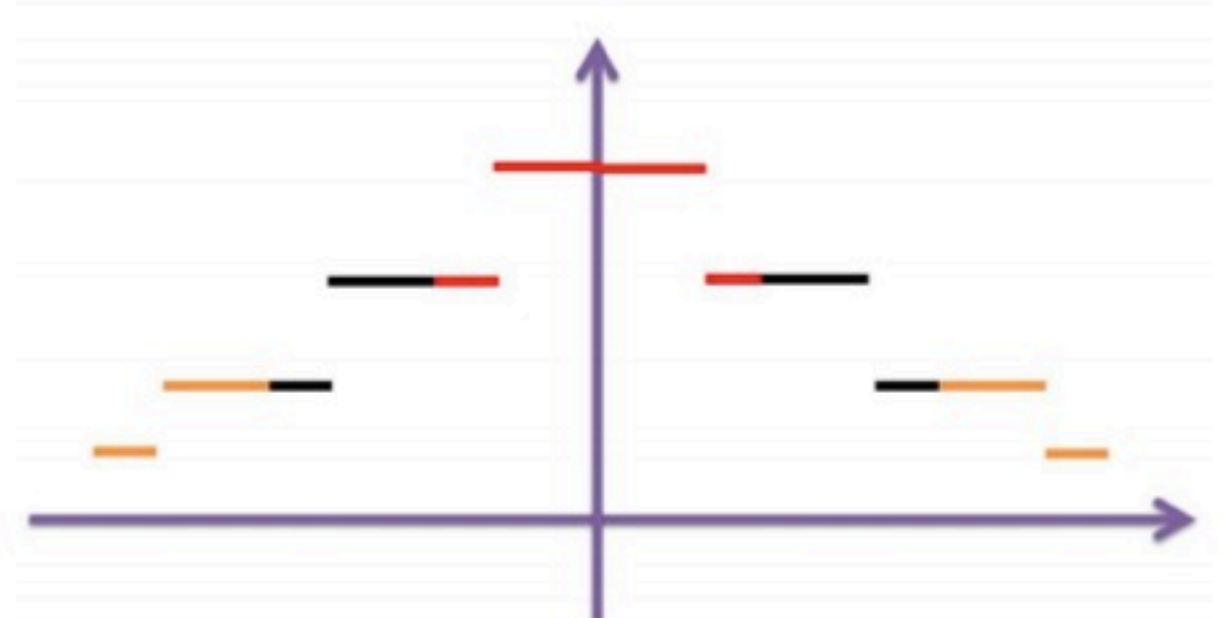
- $\min E[L(X)]$

Universal Optimality

- **Theorem: Optimal Noise is Staircase shaped**



(a) Laplace Mechanism

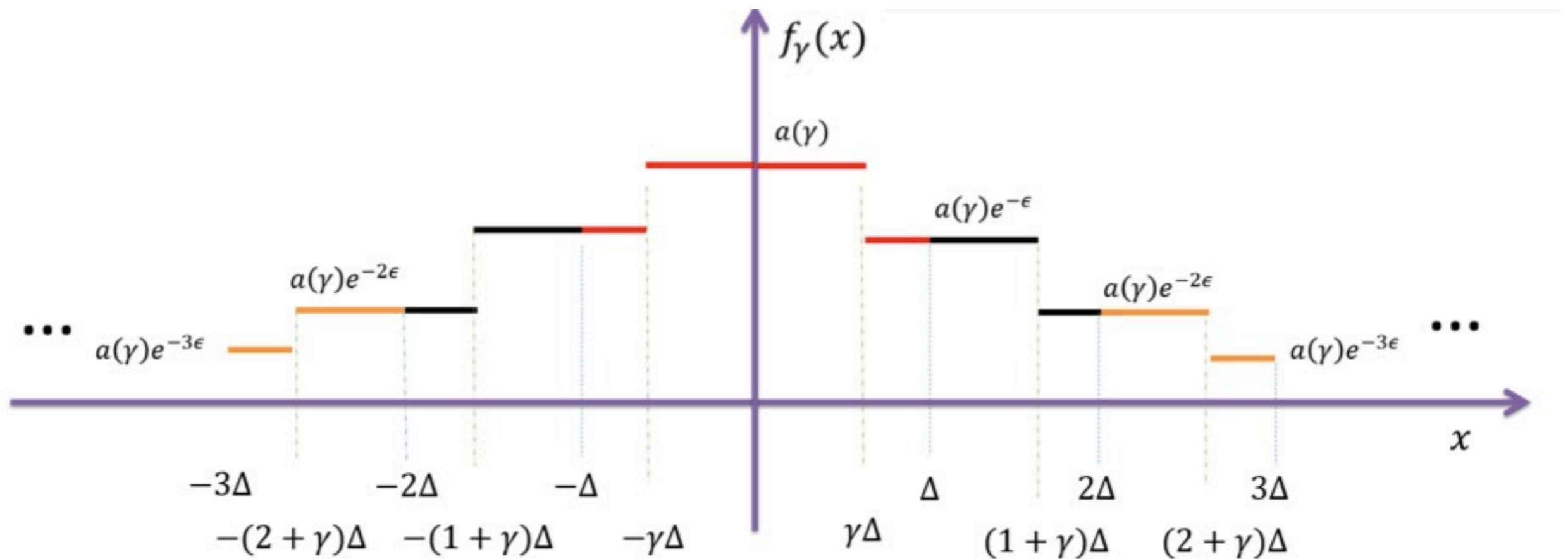


(b) Staircase Mechanism

- **Geometric mixture of uniform random variables**

Staircase Mechanism

- **Theorem:** Optimal Noise is **universally Staircase shaped**



- **Geometric decaying**
 - $\gamma \in [0, 1]$ depends on cost function

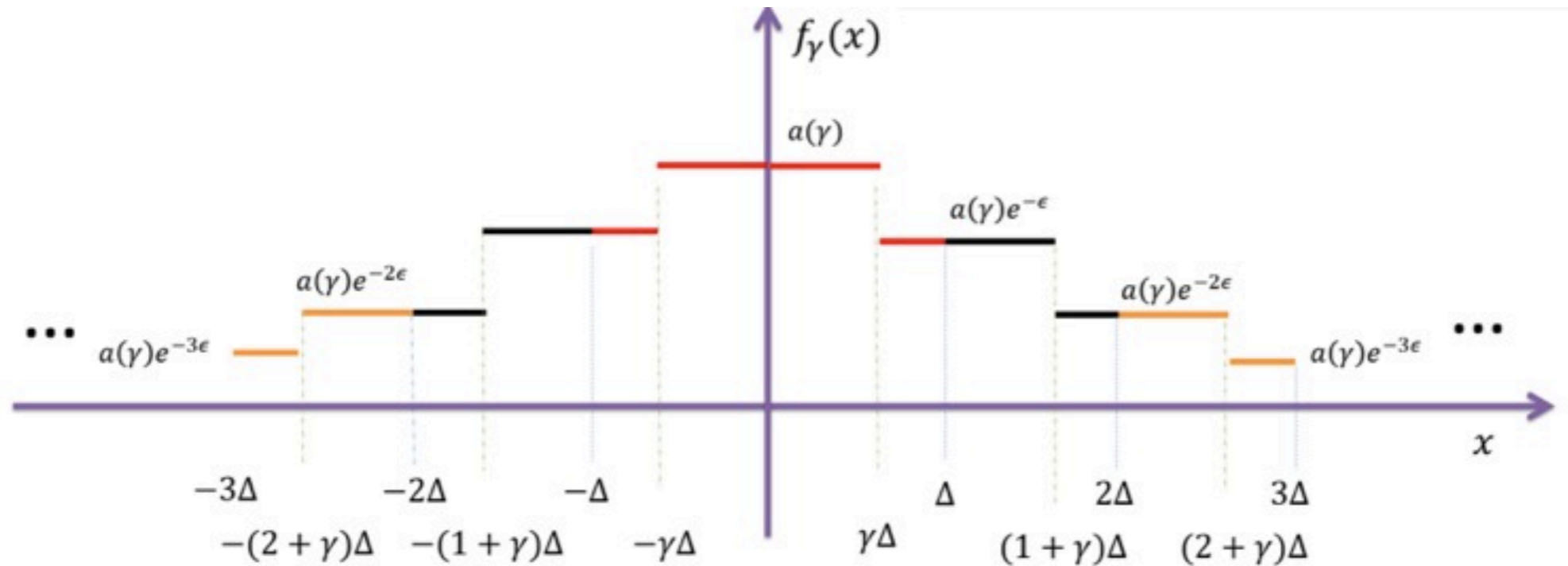
Price of Privacy

- For $L(x) = |x|$
- Minimum noise magnitude $\frac{\Delta e^{-\varepsilon/2}}{1 - e^{-\varepsilon/2}}$
- Laplace noise magnitude $\frac{\Delta}{\varepsilon}$
- High privacy
 - gap is small
- Low privacy
 - exponential improvement
- Low privacy costs exponentially less

Price of Privacy

- For $L(x) = x^2$
- Minimum noise variance $\Theta\left(\frac{\Delta^2 e^{-2\varepsilon/3}}{(1-e^{-\varepsilon})^2}\right)$
- Laplace noise variance $\frac{\Delta^2}{\varepsilon^2}$
- High privacy
 - gap is small
- Low privacy
 - exponential improvement
- Low privacy costs exponentially less

Properties of γ^*



- Need to pick γ^* ; depends on cost function

- General Properties:

$$\gamma^* \rightarrow \frac{1}{2} \quad \epsilon \rightarrow 0$$

$$\gamma^* \rightarrow 0 \quad \epsilon \rightarrow \infty$$

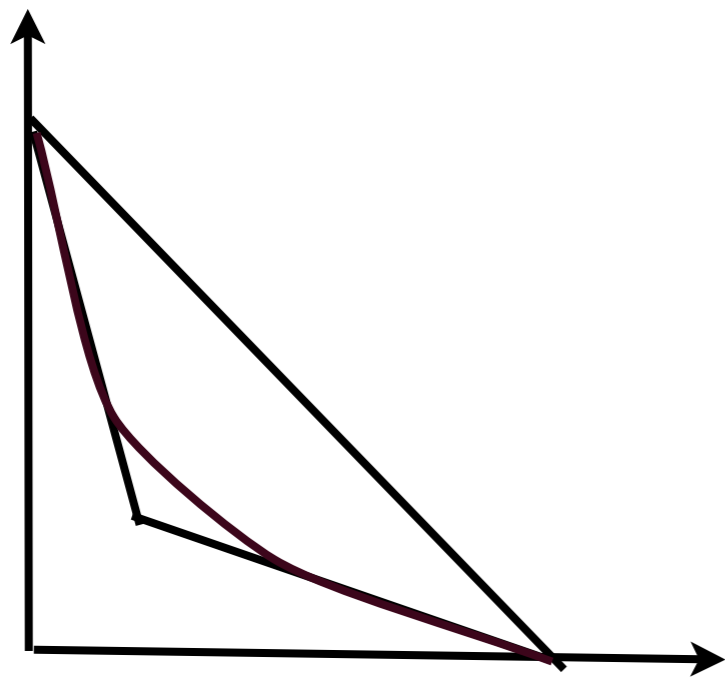
- Log Likelihood ratio $\in \{-\epsilon, 0, +\epsilon\}$

Canonical Result

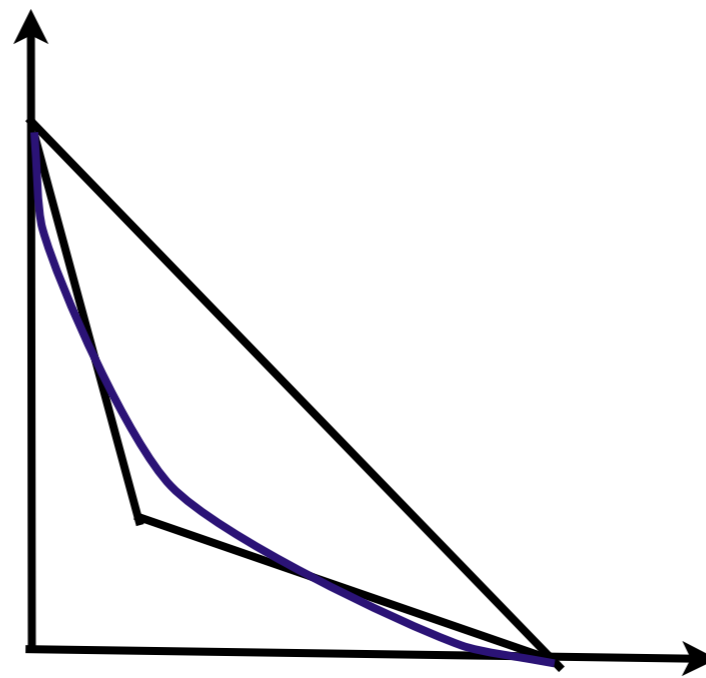
- Laplacian mechanism (and variants) widely used
 - many papers on differential privacy
- Staircase mechanism applies
 - in nearly each case
 - improves performance nearly each time
 - pronounced improvement in moderate/low privacy regimes
- Two limitations
 - intuition missing
 - generalization hard
 - data/query dependent mechanisms

FA-MD Tradeoff Curves

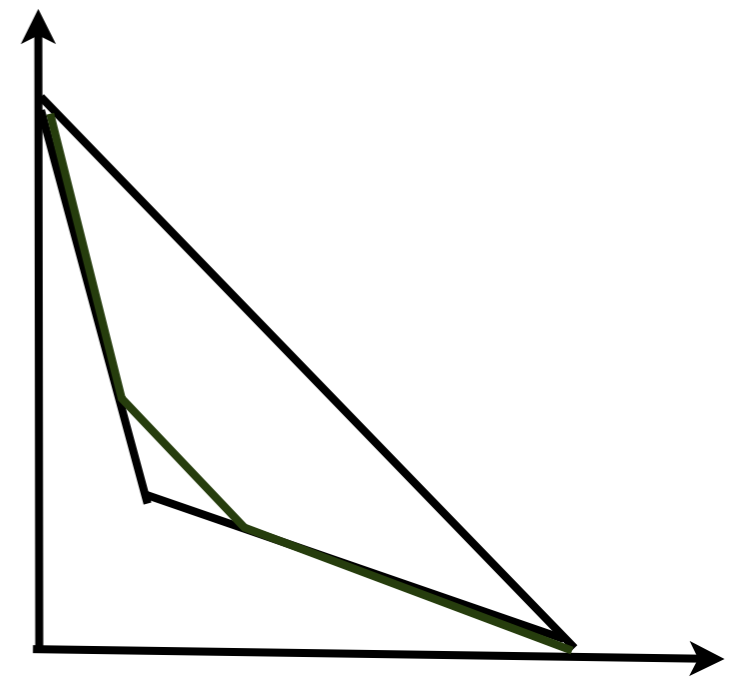
- Operational setting
- binary hypothesis testing



Laplacian



Gaussian

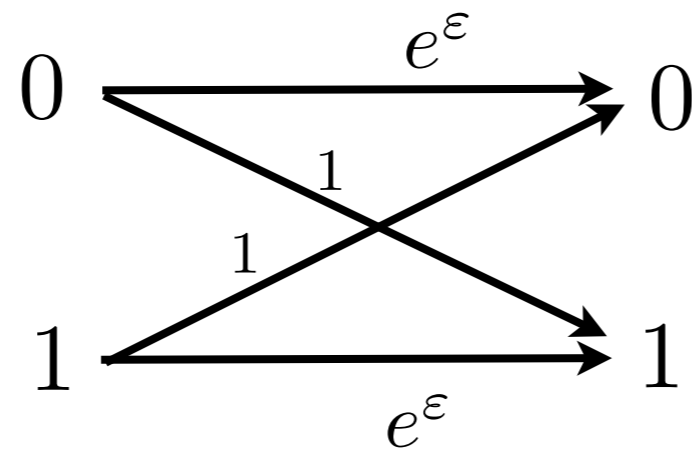


Staircase

- too complicated
- multiple query output values

Binary Query

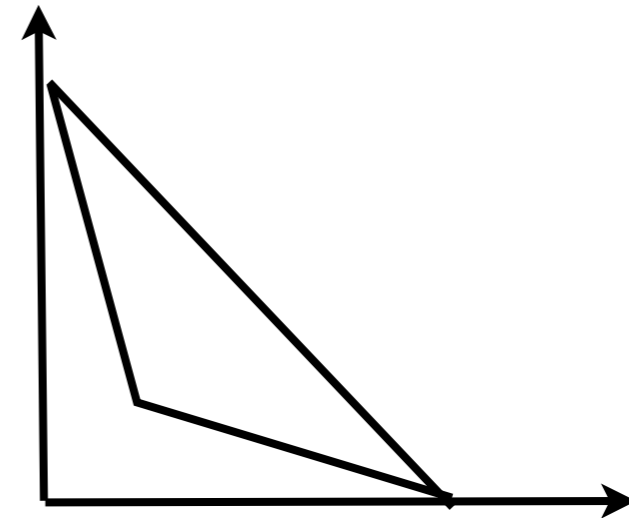
- **Binary** output
 - Yes or No answer
- Natural mechanism
 - **randomized response; W59**



- **Potentially suboptimal** in general
 - more complicated outputs
 - 2-party distributed AND computation GMPS13

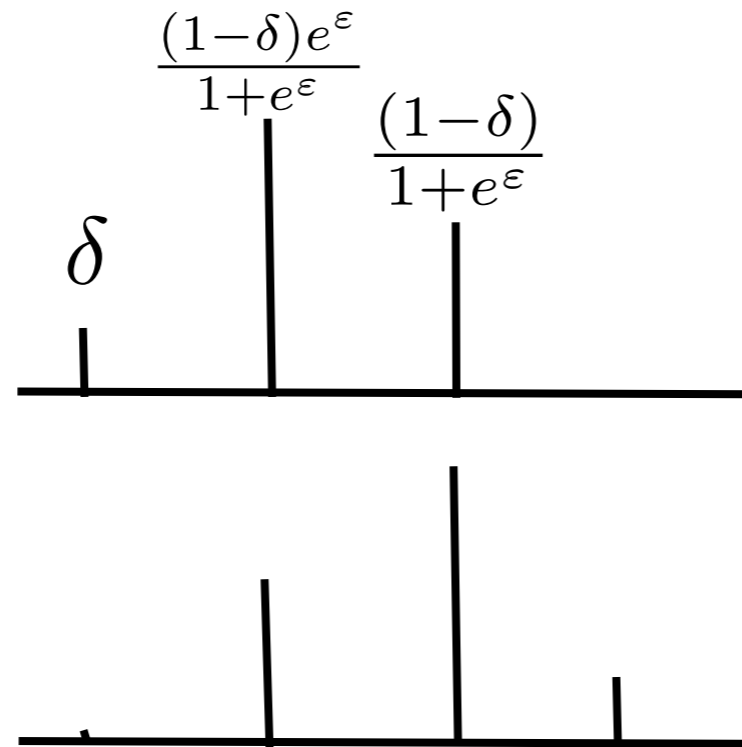
Operational Look

- **Binary** output
 - randomized response X
 - likelihood ratio $\in \{-\varepsilon, +\varepsilon\}$
- **Exactly meets** the privacy region
- Any other mechanism Y
 - only **inside the triangular region**
- Reverse Data Processing Theorem: B53
 - $D - X - Y$ -- Y can be simulated from X
 - Implications for GMPS13 -- distributed AND computation

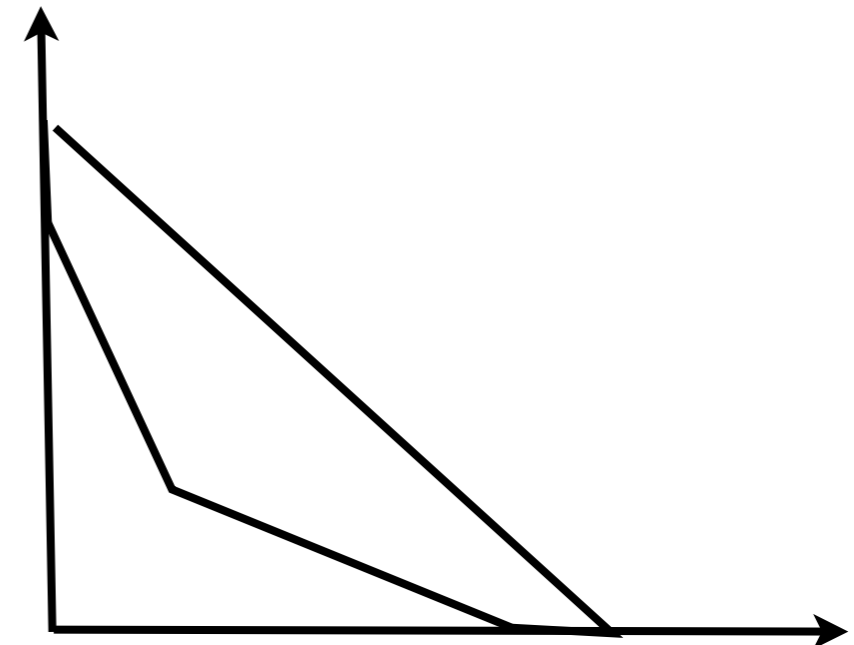


Approximate Differential Privacy

- Privatized response has **four** output letters

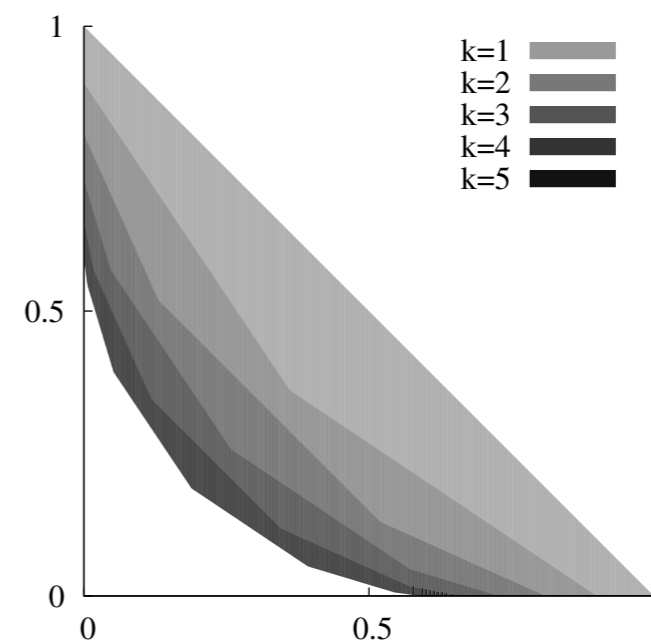
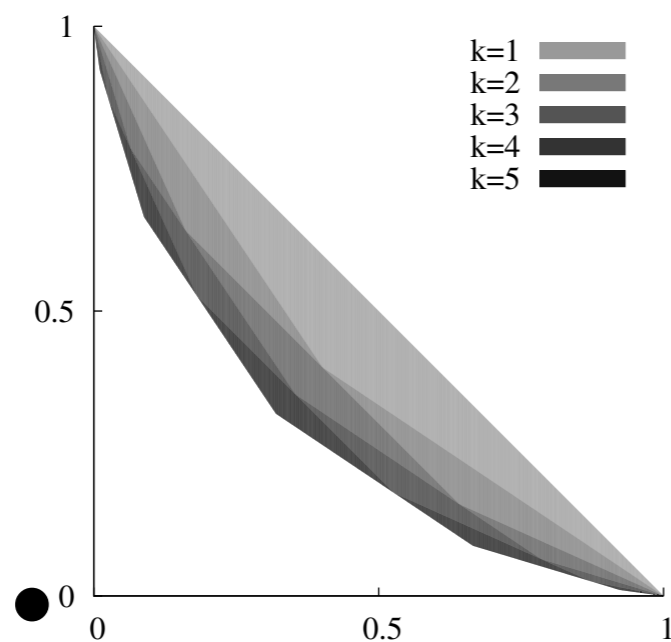


- Exactly meets the privacy region
- Any other mechanism Y
- only inside the privacy region
- $D - X - Y$



Composition Theorem

- Privacy region met exactly
- every other mechanism can be simulated
- Optimal Composition Theorem
- Composing k queries
- privacy region is intersection
- of $((k - 2i)\varepsilon, \delta_i)$ privacy regions for $i=1..k$



Composition Theorem Simplified

- **Optimal Composition Theorem**
 - conceptually straightforward
- Can be expressed as $(\tilde{\epsilon}, \delta)$ privacy
- k -fold composition, each $(\epsilon, 0)$ private

$$\tilde{\epsilon} \approx k\epsilon^2 + \epsilon \sqrt{2k \log(e + (\sqrt{k\epsilon^2}/\delta))}$$

- contrast with state of the art [DRV10]

$$\tilde{\epsilon} \approx k\epsilon^2 + \epsilon \sqrt{2k \log(1/\delta)}$$

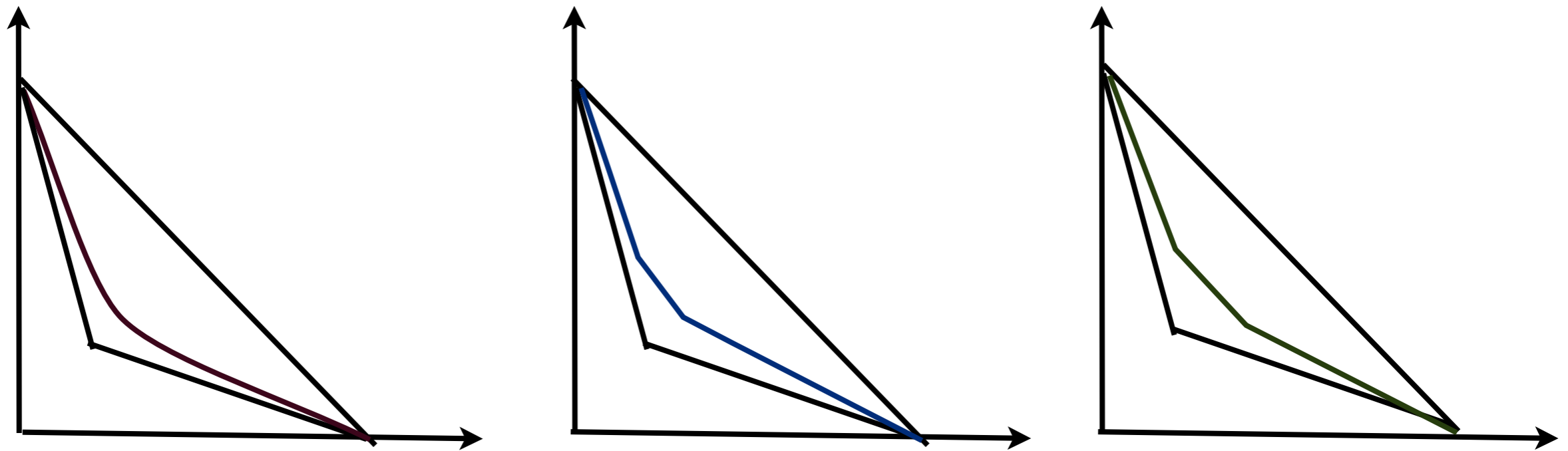
- **saving of log factor**

Applications of the Composition Theorem

- **Order optimality**
 - for many mechanisms
 - Laplace
 - Staircase
 - Gaussian
- Direct composition improves performance of Gaussian mechanism
 - sharper concentration analysis
 - chernoff bound
 - direct expression for privacy region
- **Immediate applications**
 - each intermediate step has less noise

Back to the Staircase Mechanism

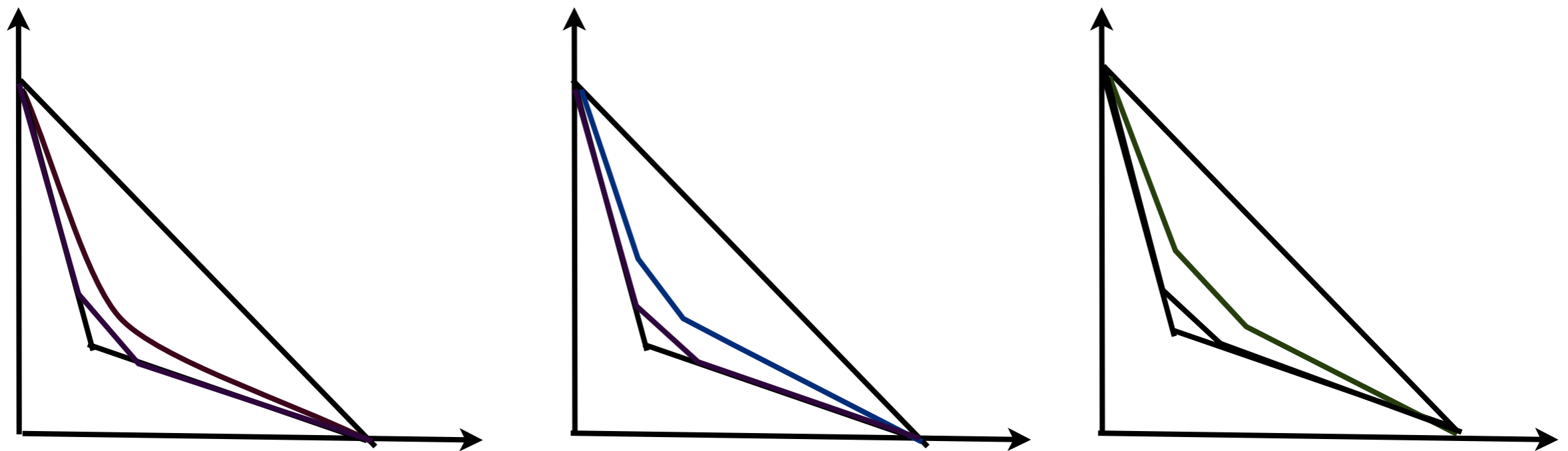
- Ternary query output
 - each pair is neighboring
- View through the **operational lens**
 - three FA-MD diagrams, one for each pair



- tradeoff among the privacy regions
 - all three regions cannot meet the full triangular region

Back to the Staircase Mechanism

- Ternary query output
 - each pair is neighboring
- Tradeoff among the privacy regions



- Staircase mechanism universally dominates
- Theorem: Every mechanism can be simulated from the staircase mechanism
- Special reverse data processing inequality

Summary

- **Fundamental Mechanisms**
 - Staircase mechanism
- **Universality**
 - cost framework
 - Markov chain framework
- **Operational Lens**
 - data processing inequalities
- **Connections to statistics**
 - Blackwell, LeCam
 - converse results to Neyman-Pearson

- Q. Geng and P. Viswanath,
- The Optimal Mechanism for Differential Privacy
- [arxiv.org/1212.1186](https://arxiv.org/abs/1212.1186)



- Q. Geng and P. Viswanath,
- The Optimal Mechanism for Differential Privacy: Multidimensional Setting
- [arxiv.org/1312.0655](https://arxiv.org/abs/1312.0655)

- S. Oh and P. Viswanath
- The Composition Theorem for Differential Privacy
- [arxiv.org/1311.0776](https://arxiv.org/abs/1311.0776)



- Q. Geng and P. Viswanath
- Optimal Mechanisms for Approximate Differential Privacy
- [arxiv.org/1305.1330](https://arxiv.org/abs/1305.1330)

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