Matrix Martingales in Randomized Numerical Linear Algebra

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Concentration of Scalar Random Variables

Random $X = \sum_{i} X_{i}$, $X_{i} \in \mathbb{R}$ 1. X_{i} are independent 2. $\mathbb{E}X_{i} = 0$

Is $X \approx 0$ with high probability?

Concentration of Scalar Random Variables

Bernstein's Inequality

Random
$$X = \sum_{i} X_{i}$$
, $X_{i} \in \mathbb{R}$
1. X_{i} are independent
2. $\mathbb{E}X_{i} = 0$
3. $|X_{i}| \leq r$
4. $\sum_{i} \mathbb{E}X_{i}^{2} \leq \sigma^{2}$

gives E.g. if $\varepsilon = 0.5$, and $r, \sigma^2 = 0.1/\log(1/\tau)$

$$\mathbb{P}[|X| > \varepsilon] \le 2\exp\left(-\frac{\varepsilon^2/2}{r\varepsilon + \sigma^2}\right)$$

Concentration of Scalar Martingales

Random $X = \sum_{i} X_{i}$, $X_{i} \in \mathbb{R}$ 1. $\frac{X_{i}}{X_{i}}$ are independent 2. $\frac{\mathbb{E}X_{i}}{\mathbb{E}X_{i}} = 0$ $\mathbb{E}[X_{i}|X_{1}, \dots, X_{i-1}] = 0$ Concentration of Scalar Martingales

Random $X = \sum_{i} X_{i}$, $X_{i} \in \mathbb{R}$ 1. $\frac{X_{i} \text{ are independent}}{\mathbb{E}[X_{i}| \text{ previous steps}] = 0$

Concentration of Scalar Martingales **Freedman's Inequality Bernstein's Inequality** Random $X = \sum_i X_i, \quad X_i \in \mathbb{R}$ 1. X_i are independent 2. $\mathbb{E}X_i = 0$ $\mathbb{E}[X_i | \text{ previous steps}] = 0$ 3. $|X_i| \le r$ 4. $\sum_{i} \mathbb{E}X_{i}^{2} \leq \sigma^{2} \quad \sum_{i} \mathbb{E}[X_{i}^{2} | \text{ previous steps}] \leq \sigma^{2}$ $\mathbb{P}\left[\sum_{i} \mathbb{E}[X_{i}^{2} | \text{ previous steps}] > \sigma^{2}\right] \leq \delta$ gives

$$\mathbb{P}[|X| > \varepsilon] \le 2 \exp\left(-\frac{\varepsilon^2/2}{r\varepsilon + \sigma^2}\right) + \delta$$

Concentration of Scalar Martingales

Freedman's Inequality

Random $X = \sum_i X_i$, $X_i \in \mathbb{R}$

- 1. X_i are independent
- 2. $\mathbb{E}[X_i | \text{ previous steps}] = 0$
- 3. $|X_i| \leq r$
- 4. $\mathbb{P}[\sum_{i} \mathbb{E}[X_{i}^{2} | \text{ previous steps}] > \sigma^{2}] \leq \delta$

gives

$$\mathbb{P}[|X| > \varepsilon] \leq 2\exp\left(-\frac{\varepsilon^2/2}{r\varepsilon + \sigma^2}\right) + \delta$$

Concentration of Matrix Random Variables

Matrix Bernstein's Inequality (Tropp 11) Deciders $\mathbf{V} = \sum_{n=1}^{\infty} \mathbf{V} = \sum_{n=1}^{\infty} \frac{d \times d}{d}$ success the

Random $X = \sum_i X_i$ $X_i \in \mathbb{R}^{d \times d}$, symmetric

- 1. X_i are independent
- 2. $\mathbb{E}X_i = 0$
- $3. \|X_i\| \le r$
- 4. $\left\|\sum_{i} \mathbb{E} X_{i}^{2}\right\| \leq \sigma^{2}$

gives

$$\mathbb{P}[\|\boldsymbol{X}\| > \epsilon] \le \frac{d}{2} \exp\left(-\frac{\varepsilon^2/2}{r\varepsilon + \sigma^2}\right)$$

Concentration of Matrix Martingales

Matrix Freedman's Inequality (Tropp 11) Random $X = \sum_i X_i$ $X_i \in \mathbb{R}^{d \times d}$, symmetric 1. X_i are independent 2. $\mathbb{E}X_i = 0$ $\mathbb{E}[X_i | \text{ previous steps}] = 0$ 3. $\|X_i\| \leq r$ 4. $\frac{\|\sum_{i} \mathbb{E} X_{i}^{2}\|}{\leq \sigma^{2}} \mathbb{P}[\|\sum_{i} \mathbb{E} [X_{i}^{2}| \text{ prev. steps}]\| > \sigma^{2}] \leq \delta$ E.g. if $\varepsilon = 0.5$, and $r, \sigma^2 = 0.1/\log(d/\tau)$ gives

$$\mathbb{P}[\|\mathbf{X}\| > \varepsilon] \le d 2\exp\left(-\frac{\varepsilon^2/2}{r\varepsilon + \sigma^2}\right) + \delta$$

 $\leq \tau + o$

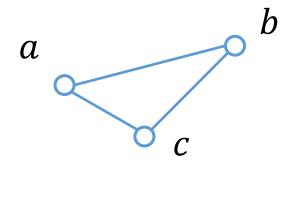
Concentration of Matrix Martingales

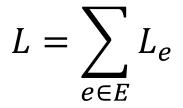
Matrix Freedman's Inequality (Tropp 11)

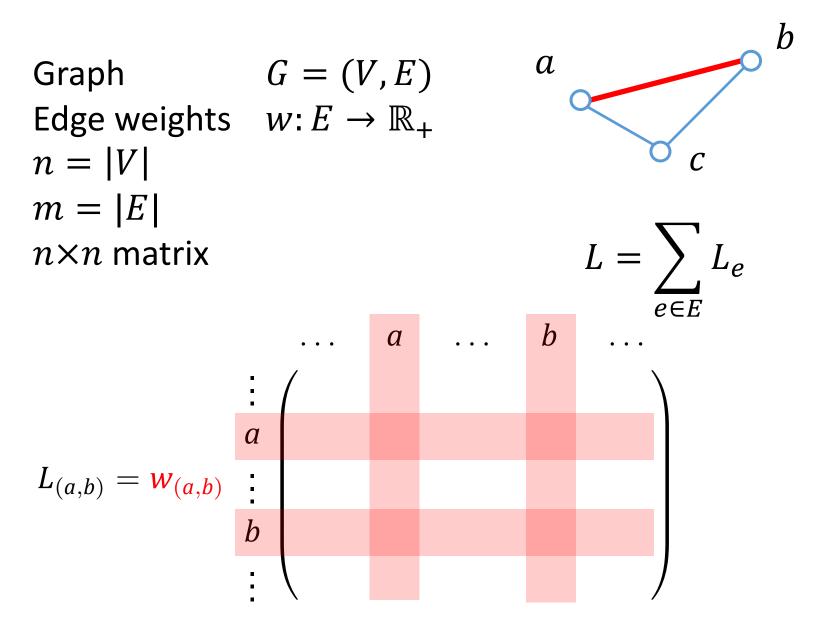
$$\mathbb{P}[\left\|\sum_{i} \mathbb{E}[X_{i}^{2} | \text{ prev. steps}]\right\| > \sigma^{2}] \leq \delta$$

Predictable quadratic variation = $\sum_{i} \mathbb{E}[X_{i}^{2}| \text{ prev. steps}]$

Graph G = (V, E)Edge weights $w: E \rightarrow \mathbb{R}_+$ n = |V|m = |E| $n \times n$ matrix







Graph G = (V, E)Edge weights $w: E \rightarrow \mathbb{R}_+$ n = |V|m = |E|

A; weighted adjacency matrix of the graph

D_{ii}dia D nali matrix of weighted degrees

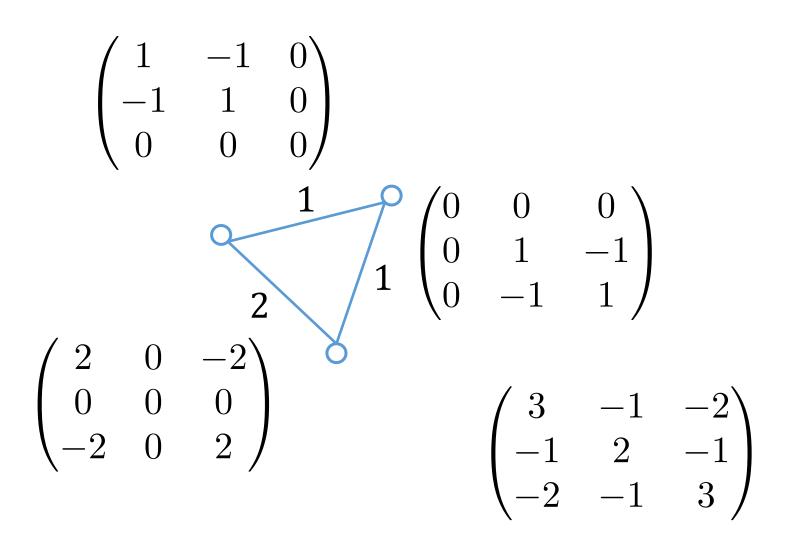
L = D - A

Symmetric matrix *L*

All off-diagonals are non-positive and

$$L_{ii} = \sum_{j \neq i} \left| L_{ij} \right|$$

Laplacian of a Graph



[ST04]: solving Laplacian linear equations in $\tilde{O}(m)$ time

[KS16]: simple algorithm

•

Solving a Laplacian Linear Equation

Lx = b

Gaussian Elimination

Find \boldsymbol{U} , upper triangular matrix, s.t. $\boldsymbol{U}^{\top}\boldsymbol{U} = \boldsymbol{L}$

Then

$$\boldsymbol{x} = \boldsymbol{U}^{-1}\boldsymbol{U}^{-\top}\boldsymbol{b}$$

Easy to apply U^{-1} and $U^{-\top}$

Solving a Laplacian Linear Equation

Lx = b

Approximate Gaussian Elimination Find U, upper triangular matrix, s.t. $U^{\top}U \approx L$

U is sparse.

 $O\left(\log\frac{1}{\varepsilon}\right)$ iterations to get ε -approximate solution \widetilde{x} .

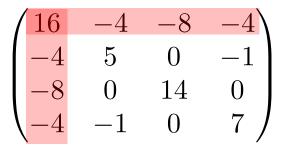
Approximate Gaussian Elimination

Theorem [KS] When L is an Laplacian matrix with m non-zeros, we can find in $O(m \log^3 n)$ time an upper triangular matrix U with $O(m \log^3 n)$ non-zeros, s.t. w.h.p.

 $\boldsymbol{U}^\top\boldsymbol{U}\approx\boldsymbol{L}$

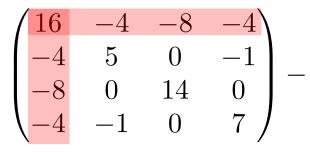
Find U, upper triangular matrix, s.t $U^{\top}U = M$

$$\boldsymbol{M} = \begin{pmatrix} 16 & -4 & -8 & -4 \\ -4 & 5 & 0 & -1 \\ -8 & 0 & 14 & 0 \\ -4 & -1 & 0 & 7 \end{pmatrix}$$



Find the rank-1 matrix that agrees with *M* on the first row and column.

$$\begin{pmatrix} 16 & -4 & -8 & -4 \\ -4 & 1 & 2 & 1 \\ -8 & 2 & 4 & 2 \\ -4 & 1 & 2 & 1 \end{pmatrix} = \begin{pmatrix} 4 \\ -1 \\ -2 \\ -1 \end{pmatrix} \begin{pmatrix} 4 \\ -1 \\ -2 \\ -1 \end{pmatrix}^{\top}$$



Subtract the rank 1 matrix. We have **eliminated the first variable.**

$$\begin{pmatrix} 16 & -4 & -8 & -4 \\ -4 & 1 & 2 & 1 \\ -8 & 2 & 4 & 2 \\ -4 & 1 & 2 & 1 \end{pmatrix}$$

$$\begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 4 & -2 & -2 \\ 0 & -2 & 10 & -2 \\ 0 & -2 & -2 & 6 \end{pmatrix}$$

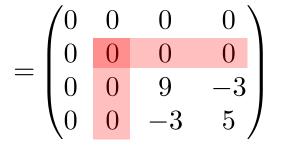
The remaining matrix is PSD.

$$\begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 4 & -2 & -2 \\ 0 & -2 & 10 & -2 \\ 0 & -2 & -2 & 6 \end{pmatrix}$$

Find rank-1 matrix that agrees with our matrix on the **next** row and column.

$$\begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 4 & -2 & -2 \\ 0 & -2 & 1 & 1 \\ 0 & -2 & 1 & 1 \end{pmatrix} = \begin{pmatrix} 0 \\ 2 \\ -1 \\ -1 \end{pmatrix} \begin{pmatrix} 0 \\ 2 \\ -1 \\ -1 \end{pmatrix}^{\top}$$

$$\begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 4 & -2 & -2 \\ 0 & -2 & 10 & -2 \\ 0 & -2 & -2 & 6 \end{pmatrix}$$



Subtract the rank 1 matrix. We have **eliminated the second variable.**

$$\begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 4 & -2 & -2 \\ 0 & -2 & 1 & 1 \\ 0 & -2 & 1 & 1 \end{pmatrix}$$

Repeat until all parts written as rank 1 terms.

$$\mathbf{M} = \begin{pmatrix} 16 & -4 & -8 & -4 \\ -4 & 5 & 0 & -1 \\ -8 & 0 & 14 & 0 \\ -4 & -1 & 0 & 7 \end{pmatrix}$$
$$= \begin{pmatrix} 4 \\ -1 \\ -2 \\ -1 \end{pmatrix} \begin{pmatrix} 4 \\ -1 \\ -2 \\ -1 \end{pmatrix}^{\top} + \begin{pmatrix} 0 \\ 2 \\ -1 \\ -1 \end{pmatrix} \begin{pmatrix} 0 \\ 2 \\ -1 \\ -1 \end{pmatrix}^{\top} + \begin{pmatrix} 0 \\ 0 \\ 3 \\ -1 \end{pmatrix} \begin{pmatrix} 0 \\ 0 \\ 3 \\ -1 \end{pmatrix}^{\top} + \begin{pmatrix} 0 \\ 0 \\ 0 \\ 2 \\ 2 \end{pmatrix} \begin{pmatrix} 0 \\ 0 \\ 0 \\ 2 \end{pmatrix}^{\top}$$

Repeat until all parts written as rank 1 terms.

$$\boldsymbol{M} = \begin{pmatrix} 16 & -4 & -8 & -4 \\ -4 & 5 & 0 & -1 \\ -8 & 0 & 14 & 0 \\ -4 & -1 & 0 & 7 \end{pmatrix}$$
$$= \begin{pmatrix} 4 & 0 & 0 & 0 \\ -1 & 2 & 0 & 0 \\ -2 & -1 & 3 & 0 \\ -1 & -1 & -1 & 2 \end{pmatrix} \begin{pmatrix} 4 & -1 & -2 & -1 \\ 0 & 2 & -1 & -1 \\ 0 & 0 & 3 & -1 \\ 0 & 0 & 0 & 2 \end{pmatrix}$$

Repeat until all parts written as rank 1 terms.

$$\begin{split} \boldsymbol{M} &= \begin{pmatrix} 16 & -4 & -8 & -4 \\ -4 & 5 & 0 & -1 \\ -8 & 0 & 14 & 0 \\ -4 & -1 & 0 & 7 \end{pmatrix}^{\top} \\ &= \begin{pmatrix} 4 & -1 & -2 & -1 \\ 0 & 2 & -1 & -1 \\ 0 & 0 & 3 & -1 \\ 0 & 0 & 0 & 2 \end{pmatrix}^{\top} \begin{pmatrix} 4 & -1 & -2 & -1 \\ 0 & 2 & -1 & -1 \\ 0 & 0 & 3 & -1 \\ 0 & 0 & 0 & 2 \end{pmatrix} = \boldsymbol{U}^{\top} \boldsymbol{U} \end{split}$$

What is special about Gaussian Elimination on Laplacians?

The **remaining matrix** is always Laplacian.

$$\boldsymbol{L} = \begin{pmatrix} 16 & -8 & -4 & -4 \\ -8 & 8 & 0 & 0 \\ -4 & 0 & 4 & 0 \\ -4 & 0 & 0 & 4 \end{pmatrix}$$

What is special about Gaussian Elimination on Laplacians?

The remaining matrix is always Laplacian.

$$\boldsymbol{L} = \begin{pmatrix} 16 & -8 & -4 & -4 \\ -8 & 4 & 2 & 2 \\ -4 & 2 & 1 & 1 \\ -4 & 2 & 1 & 1 \end{pmatrix} + \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & 4 & -2 & -2 \\ 0 & -2 & 3 & -1 \\ 0 & -2 & -1 & 3 \end{pmatrix}$$

What is special about Gaussian Elimination on Laplacians?

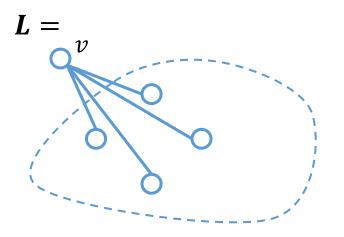
The remaining matrix is always Laplacian.

$$\boldsymbol{L} = \begin{pmatrix} 4 \\ -2 \\ -1 \\ -1 \\ -1 \end{pmatrix} \begin{pmatrix} 4 \\ -2 \\ -1 \\ -1 \end{pmatrix}^{\top} + \begin{pmatrix} 0 & 0 & 0 & 0 \\ 0 & \boldsymbol{4} & -2 & -2 \\ 0 & -2 & \boldsymbol{3} & -1 \\ 0 & -2 & -1 & \boldsymbol{3} \end{pmatrix}$$

A new Laplacian!

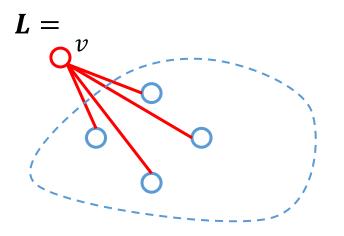
Solving Lx = b by Gaussian Elimination can take $\Omega(n^3)$ time.

The main issue is **fill**



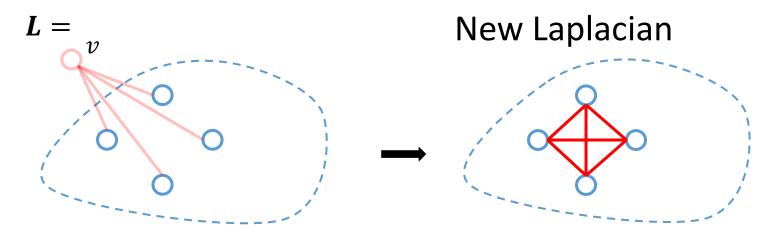
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Solving Lx = b by Gaussian Elimination can take $\Omega(n^3)$ time.

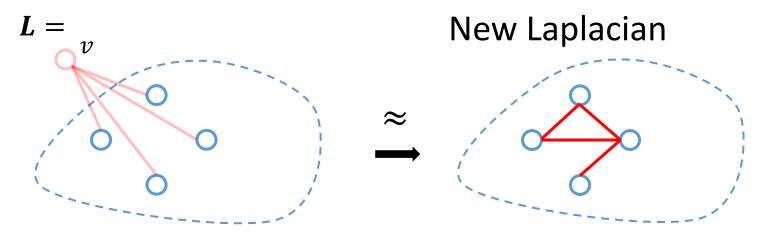
The main issue is **fill**



Elimination creates a clique on the neighbors of v

Solving Lx = b by Gaussian Elimination can take $\Omega(n^3)$ time.

The main issue is **fill**



Laplacian cliques can be sparsified!

Gaussian Elimination

- 1. Pick a vertex v to eliminate
- 2. Add the clique created by eliminating v
- 3. Repeat until done

Approximate Gaussian Elimination

- 1. Pick a vertex v to eliminate
- 2. Add the clique created by eliminating v
- 3. Repeat until done

Approximate Gaussian Elimination

- 1. Pick a random vertex v to eliminate
- 2. Add the clique created by eliminating v
- 3. Repeat until done

Approximate Gaussian Elimination

- 1. Pick a random vertex v to eliminate
- 2. Sample the clique created by eliminating v
- 3. Repeat until done

Resembles randomized Incomplete Cholesky

Approximating Matrices by Sampling

Goal

$\boldsymbol{U}^\top \boldsymbol{U} \approx \boldsymbol{L}$

Approach

- 1. $\mathbb{E} U^{\top}U = L$
- 2. Show $\boldsymbol{U}^{\top}\boldsymbol{U}$ concentrated around expectation

Gives $U^{\top}U \approx L$ w. high probability

$$\begin{pmatrix} L^{(0)} \end{pmatrix} = \begin{pmatrix} c_1 \end{pmatrix} \begin{pmatrix} c_1 \end{pmatrix}^{\mathsf{T}} + \begin{pmatrix} \bullet & \bullet & \bullet \\ \bullet & \bullet & \bullet \\ \bullet & \bullet & \bullet \end{pmatrix}$$

Original
Laplacian term + clique

$$= \left(c_1 \right) \left(c_1 \right)^{\mathsf{T}} + \left(\begin{array}{c} \bullet \bullet \bullet \\ \bullet \bullet \bullet \bullet \\ \bullet \bullet \bullet \bullet \end{array} \right)$$

Rank 1
term Remaining graph
+ sparsified clique

$$\begin{pmatrix} L^{(1)} \end{pmatrix} = \begin{pmatrix} c_1 \end{pmatrix} \begin{pmatrix} c_1 \end{pmatrix}^{\mathsf{T}} + \begin{pmatrix} \bullet & \bullet & \bullet \\ \bullet & \bullet & \bullet \\ \bullet & \bullet & \bullet \end{pmatrix}$$

Rank 1
term + sparsified clique

$$\mathbb{E}\left(\begin{array}{c} L^{(1)} \end{array}\right) = \left(\begin{array}{c} c_1 \end{array}\right) \left(\begin{array}{c} c_1 \end{array}\right)^{\mathsf{T}} + \mathbb{E}\left(\begin{array}{c} \bullet \bullet \bullet \\ \bullet \bullet \bullet \\ \bullet \bullet \bullet \end{array}\right)$$

Rank 1
term + sparsified clique

$$\mathbb{E}\left(\begin{array}{c} L^{(1)} \end{array}\right) = \left(\begin{array}{c} c_1 \end{array}\right) \left(\begin{array}{c} c_1 \end{array}\right)^{\mathsf{T}} + \mathbb{E}\left(\begin{array}{c} \bullet \bullet \bullet \\ \bullet \bullet \bullet \\ \bullet \bullet \bullet \end{array}\right)$$
Rank 1
Remaining graph
term
$$\operatorname{Remaining graph} + \operatorname{sparsified clique}$$
Suppose
$$= \left(\begin{array}{c} c_1 \end{array}\right) \left(\begin{array}{c} c_1 \end{array}\right)^{\mathsf{T}} + \left(\begin{array}{c} \bullet \bullet \bullet \\ \bullet \bullet \bullet \\ \bullet \bullet \bullet \end{array}\right)$$

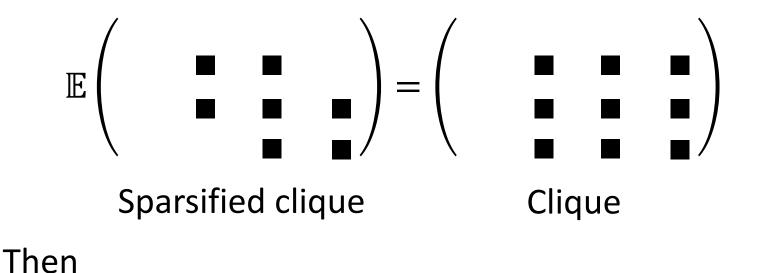
$$\mathbb{E}\left(\begin{array}{c} \boldsymbol{L}^{(1)}\end{array}\right) = \left(\boldsymbol{c}_{1}\right)\left(\boldsymbol{c}_{1}\right)^{\mathsf{T}} + \mathbb{E}\left(\begin{array}{c} \bullet \bullet \bullet \\ \bullet \bullet \bullet \\ \bullet \bullet \bullet \end{array}\right)$$

$$\begin{array}{c} \mathsf{Rank 1} \\ \mathsf{term} \end{array} \quad \begin{array}{c} \mathsf{Remaining graph} \\ \mathsf{+ sparsified clique} \end{array}$$

$$\begin{array}{c} \mathsf{Then} = \boldsymbol{L}^{(0)} \end{array}$$

Let $L^{(i)}$ be our approximation after *i* eliminations

If we ensure at each step



$$\mathbb{E} L^{(i)} = L^{(i-1)}$$
$$\mathbb{E} \left[L^{(i)} - L^{(i-1)} \right] \text{ previous steps} = 0$$

Approximation?

Approximate Gaussian Elimination Find U, upper triangular matrix, s.t. $U^{\top}U \approx L$

$$- \|\boldsymbol{U}^\top \boldsymbol{U} - \boldsymbol{L}\| \le 0.5$$

$$\left\|\boldsymbol{L}^{-1/2}\boldsymbol{U}^{\mathsf{T}}\boldsymbol{U}\boldsymbol{L}^{-1/2}-\boldsymbol{I}\right\| \leq 0.5$$

Essential Tools

Isotropic position $\overline{Z} \stackrel{\text{def}}{=} L^{-1/2} Z L^{-1/2}$

Goal is now

 $\overline{\boldsymbol{U}^{\top}\boldsymbol{U}}-\boldsymbol{I}\approx\boldsymbol{0}$

PSD Order

 $A \leq B$

iff for all \boldsymbol{x}

 $x^{\top}Ax \leq x^{\top}Bx$

Matrix Concentration: Edge Variables

$$\mathbf{Y}_{e} = \begin{cases} \frac{1}{p_{e}} \mathbf{L}_{e} & \text{w. probability } p_{e} \\ \mathbf{0} & \text{o.w.} \end{cases}$$

Zero-mean variables

$$X_e = Y_e - L_e$$

Isotropic position variables

 \overline{X}_e

Predictable Quadratic Variation

Predictable quadratic variation = $\sum_{i} \mathbb{E}[X_{i}^{2}| \text{ prev. steps}]$

Want to show

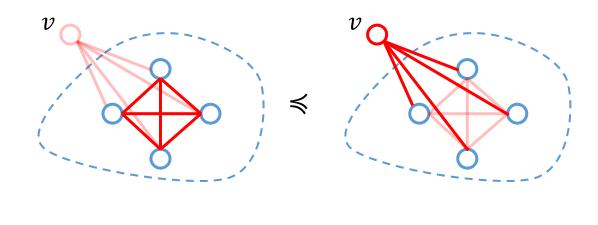
$$\mathbb{P}[\left\|\sum_{i} \mathbb{E}[X_{i}^{2} | \text{ prev. steps}]\right\| > \sigma^{2}] \leq \delta$$

Promise:

$$\mathbb{E}\overline{X}_e^2 \preccurlyeq r\overline{L}_e$$

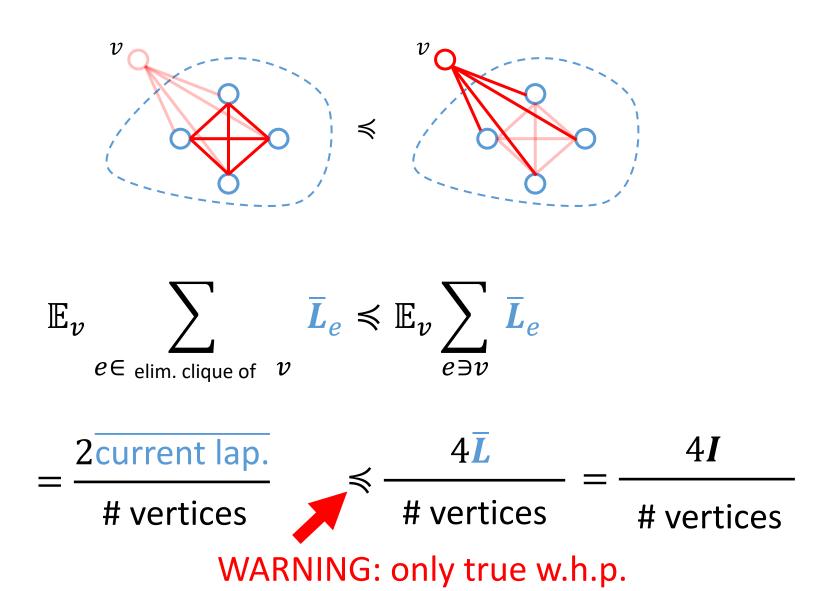


Sample Variance



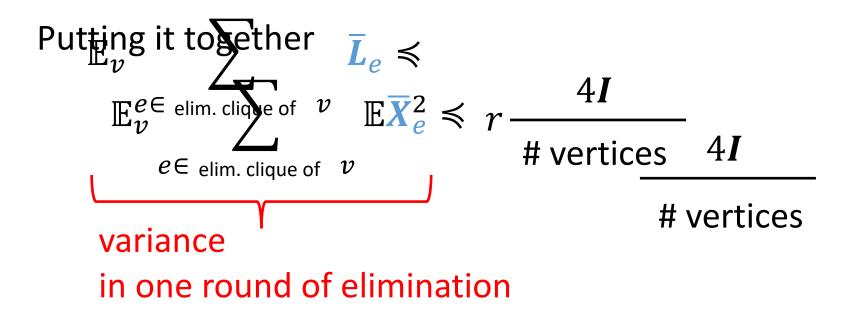
 $\bar{L}_e \preccurlyeq \sum \bar{L}_e$ $\overline{e \ni v}$ $e \in {
m elim.}$ clique of ~ v



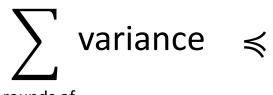


Sample Variance

Recall
$$\mathbb{E}\overline{X}_e^2 \preccurlyeq r\overline{L}_e$$

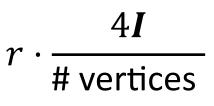


Sample Variance

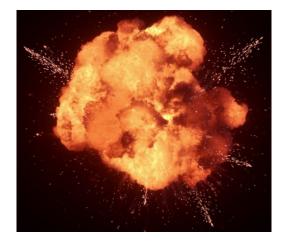


rounds of elimination





$\leq 4 r \log n \cdot I$



Summary

Matrix martingales: a natural fit for algorithmic analysis

Understanding the Predictable Quadratic Variation is key

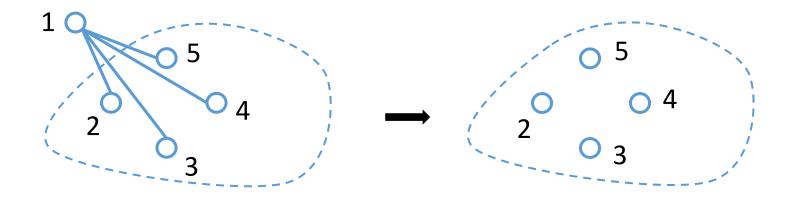
Some results using matrix martingales

Cohen, Musco, Pachocki '16 – online row sampling Kyng, Sachdeva '17 – approximate Gaussian elimination Kyng, Peng, Pachocki, Sachdeva '17 – semi-streaming graph sparsification Cohen, Kelner, Kyng, Peebles, Peng, Rao, Sidford '18 – solving Directed Laplacian eqs. Kyng, Song '18 – Matrix Chernoff bound for negatively dependent random variables

Thanks!

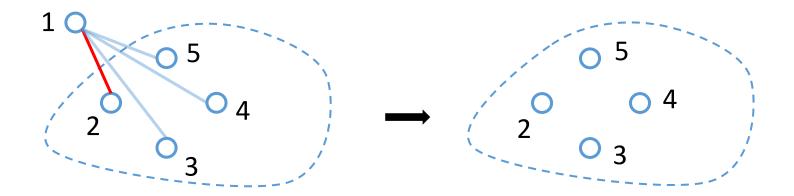
For each edge
$$(1, v)$$

pick an edge $(1, u)$ with probability $\sim w_u$
insert edge (v, u) with weight $\frac{w_u w_v}{w_u + w_v}$



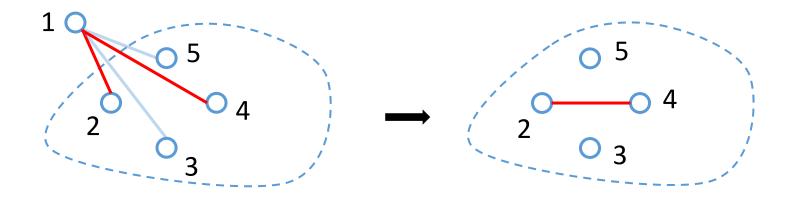
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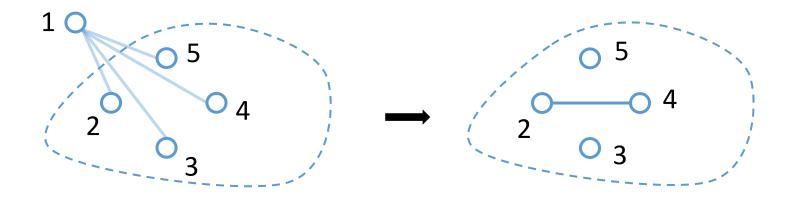
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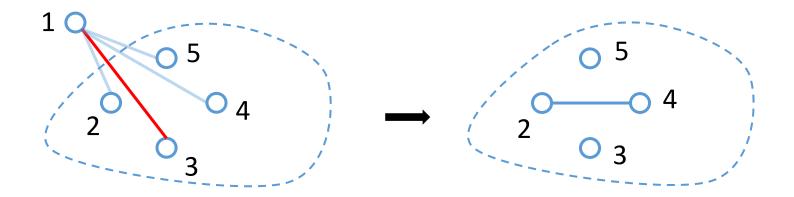
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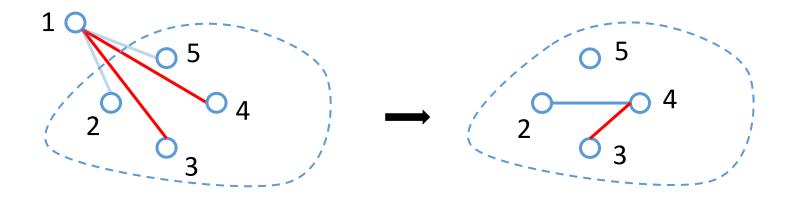
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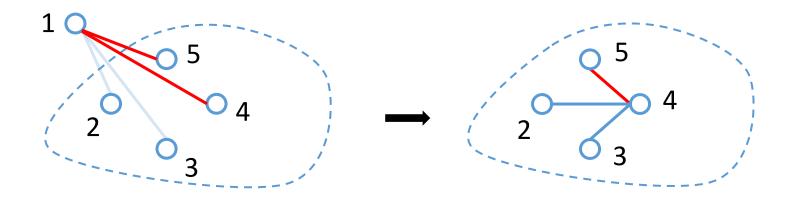
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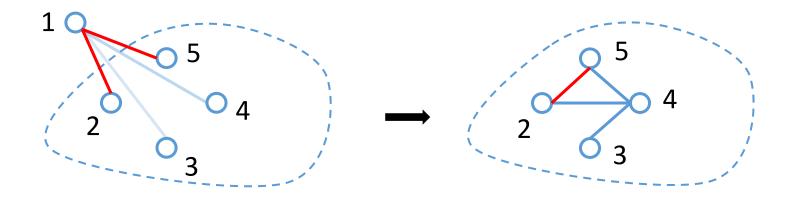
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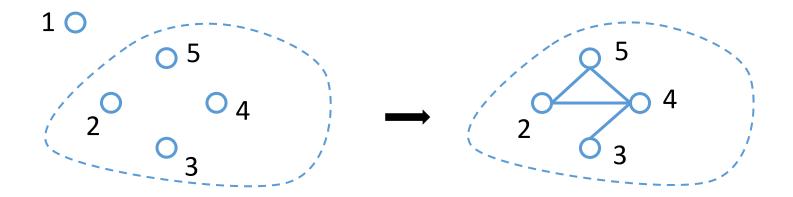
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$$(1, v)$$

pick an edge $(1, u)$ with probability $\sim w_u$
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Practice

Julia Package: Laplacians.jl

tiny.cc/spielman-solver-code

Theory

Nearly-linear time Directed Laplacian solvers using approximate Gaussian elimination [CKKPPSR17]

This is really the end